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A One Pass Algorithm for Compiling ALGOL 68 Declarations

Victor B. Schneider

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Victor B. Schneider

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A ONE-PASS ALGORITHM FOR COMPILING ALGOL 68 DECLARATIONS

I. Introduction

This report describes work underway in the design of an ALGOL 68 compiler system. It presupposes some familiarity with the general ideas of the language, such as might be gained by reading the ALGOL 68 report (1) or some of its companion documents (2,3). In what follows, we are primarily concerned with problems involved in compiling ALGOL 68 data declarations, although mention is made of techniques used in other portions of the compiler as well. Our approach is to explain the motivation behind choosing the translation grammar in Appendix 1 as a means of describing the structure and translation of programs written in the ALGOL 68 language.

The context-free grammar of Appendix 1 was liberally modeled on the Van Wijngaarden grammar of the report (1). That is to say, an attempt was made to force the rules of the grammar to resemble the rules of the report as closely as possible, except in those cases--such as the coercions of Sections 8.2.1 through 8.2.6--where the context-free grammar was clearly not adequate (and is not used) for descriptive purposes. In the case of the rules of Sections 8.2.1 through 8.2.6, the actions represented by rules in the report are to be carried out in the translation rules of the Appendix 1 grammar. Our grammar diverges from the official document principally in minor points, such as the inclusion of rules to describe the so-called "extended language" in Section 9 of the report, and of certain minor syntactic restrictions arising from the one-pass nature of our algorithm, as noted below.

The translation-grammar approach used in this report is based on earlier work by Wirth and Weber (4), Lewis and Stearns (5), Schneider (6,7), and Vere (8). The notation used in this approach can be explained in terms of the following example:

In the translation grammar of Appendix 1, the rule defining a mode declaration would be written as follows, if we could ignore the table entries and compiled code required for such a declaration:

$$\langle \text{mode declaration} \rangle \rightarrow \quad (1)$$

$$M\emptyset DE \langle \text{mode indicant} \rangle = \langle \text{actual mode declarer} \rangle$$

Rule (1) is exactly the sort of rule one would expect to find in a document such as the ALGOL 60 report (1) which contains a context-free grammar written in B.N.F. In order to describe the process of compiling translated code for a mode declaration, we add information to the context-free rule, as in Appendix 1:

$$\begin{aligned} P_{25a}: \langle \text{mode declaration} \rangle \rightarrow \\ M\emptyset DE / \psi_{25a1} / \langle \text{mode indicant} \rangle = / \psi_{25a2} / \\ \langle \text{actual mode declarer} \rangle / \psi_{25a3} / \end{aligned}$$

In this expansion of rule (1) above, we have added the notation " $P_{25a}:$ " to indicate that this rule is in the 25th group of rules in our grammar. We have also added three compile-time subroutines, namely, ψ_{25a1} , ψ_{25a2} , and ψ_{25a3} . These compile-time subroutines make entries on compile-time tables and generate code to be written out for the program scanned by the compiler. Thus, when the symbol " $M\emptyset DE$ " is recognized by the compiler, subroutine ψ_{25a1} is called; when the sequence of symbols

$$M\emptyset DE \langle \text{mode indicant} \rangle =$$

is recognized by the compiler, subroutine ψ_{25a2} is called; and

the completion of the right-hand side of syntactic rule

P_{25a} causes subroutine ψ_{25a3} to be called.

Thus, the ensemble of rules in a translation grammar may be thought of as an abstract representation of the program input to a compiler-compiler system.

In fact, our translation grammar is intended to be transferred to cards and read in to a compiler-compiler such as the one described in (6). As can be seen in Appendix 2, the subroutines for the data declaration section of our translation grammar have been programmed in FORTRAN V (no less), and it is the operation and interaction of these compiler subroutines that is the main burden of this report.

II. Techniques Used in Translation

Representation of ALGOL 68 Structures and Data Types in the Compiler

ALGOL 68 structures and invented data types are represented as linked lists in the computer memory used by the compiler. Such a representation is, naturally, machine dependent, so we propose here to give examples of how these lists are represented as storage structures on the Purdue C.D.C. 6500 computer. Since the C.D.C. 6000 series of computers has 60-bit memory words, we chose to divide each list word into three address-size fields and one character-size field for storing miscellaneous information. An example follows, showing some typical ALGOL 68 mode declarations and their compile-time representations.

Some typical mode declarations:

```

MODE $P$ = PROC(INT,$S$)REAL;
MODE $$ = STRUCT(REF $$$ POINTER, $Q$ VALUE);
MODE $Q$ = UNION($T$, $RS$);
MODE $RS$ = REF $$;
MODE $T$ = [1:5, 1:6]REF $P$;

```

Routines Used in Manipulating Compiler List Structures

The storage structures used in Figure 1 are similar to ones already suggested by Goos (9). To construct them, the compiler subroutines use a package of FORTRAN subroutines derived from the primitives of Weizenbaum's SLIP system (10). These SLIP-inspired primitives are listed and explained briefly in Table 1 below.

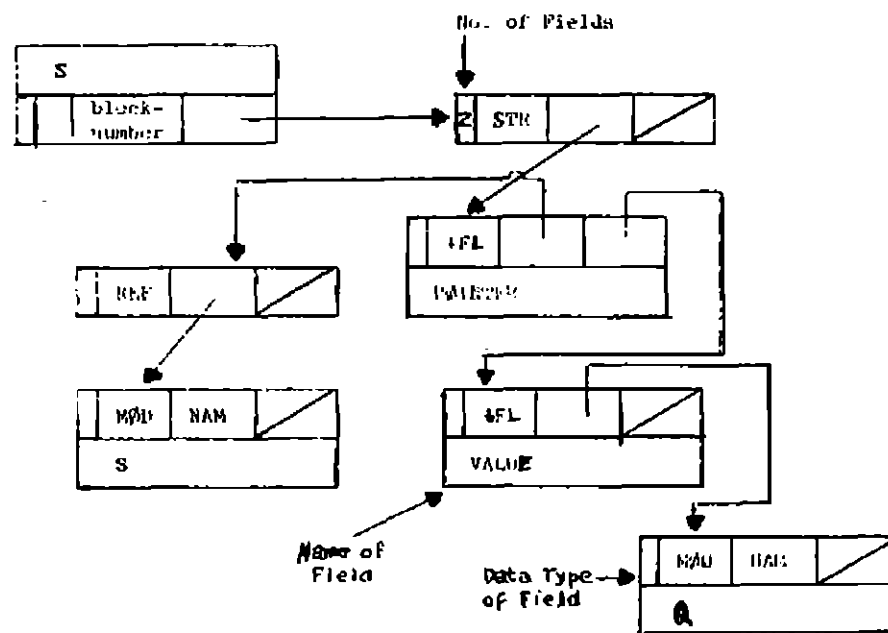
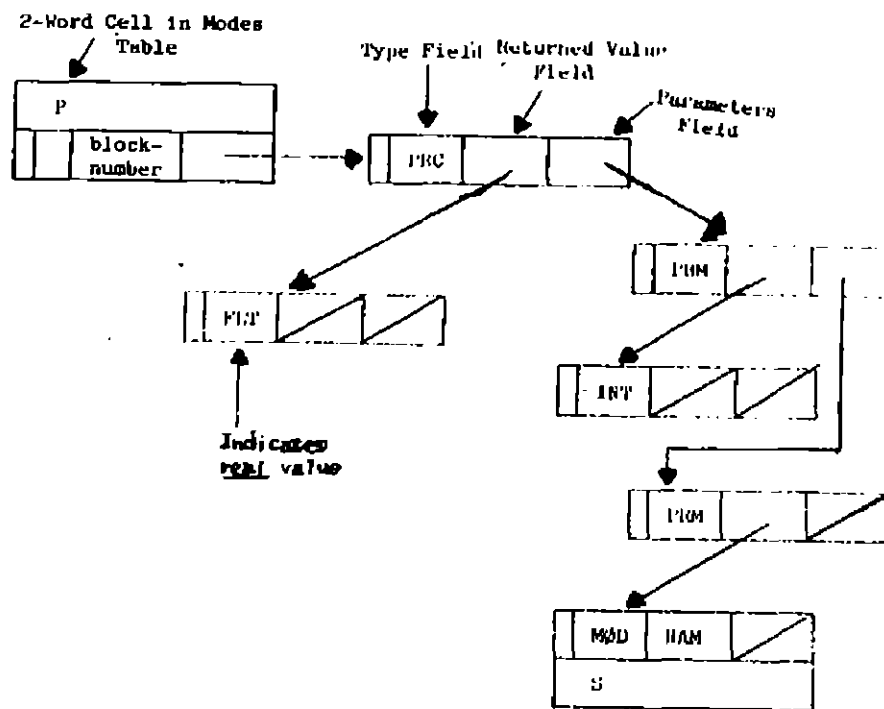


Figure 1. Representation of mode declarations by compiler lists

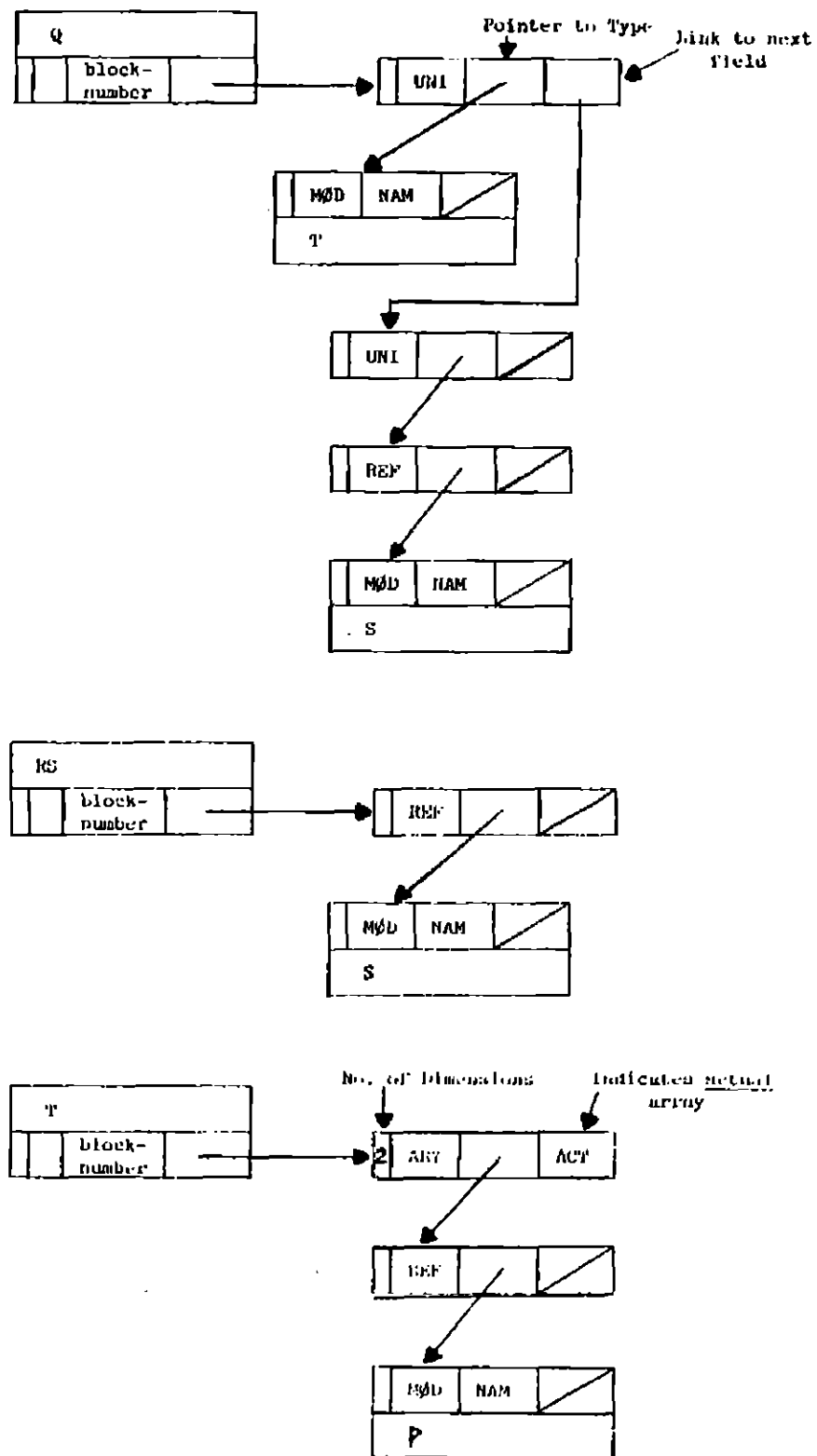


Figure 1. Continued

Table 1. List Processing Primitives Used by the Compiler

<u>Name</u>	<u>Meaning</u>
LØCF(X)	address (X) + LØCF (The function returns the machine address of FØRTRAN variable X.)
CØNT(I)	[I] → CØNT
INHALT(I)	[I] → INHALT (The value stored in the memory cell whose address is given by I is returned.)
IBYTE(J,K)	J(6*(K-1) → 6*K) → IBYTE (The Kth character in word J is returned right justified with padded zeroes.)
LINK(J,K)	J(18*(K-1) → 18*K) → LINK (The Kth address-sized field in word J is returned right justified with padded zeros.)
PUTDIR(A,B,K)	A(0 → 5) → B(6*(K-1) → 6*K) (The first character of A is put into the Kth character position of word B.)
PUTIND(A,J,K)	A(0 → 5) → [J](6*(K-1) → 6*K) (The first character of A is put into the Kth character position of the word whose machine address is stored in J.)
SETDIR(A,B,K)	A(0 → 17) → B(18*(K-1) → 18*K)
SETIND(A,J,K)	A(0 → 17) → [J](18*(K-1) → 18*K)

In addition to the functions in Table 1, there is an integer function called IFETCH(K) and a logical function called SAMETYP(I,J,K).

The purpose of IFETCH is to return the machine address of the first word in a list cell of K words stored sequentially in computer memory. The SAMETYP function is, as the name suggests, a function for comparing two compiler data structures to discover if they are of the same "type". Because of the design of ALGOL 68, two declarations may be of the same

"type" if one has any number of intermixed procedure and reference prefixes followed by a copy of the second declaration. Thus, for example, if parameter I pointed to a declaration, such as "PRØC REF REF PRØC REAL", and J pointed to "PRØC REAL", the SAMETYP function would return the FORTRAN .TRUE. value. In addition, after execution of the function, the K parameter would point to a newly-created list containing the sequence "PRØC REF REF". Thus, the list returned to K indicates how many levels of dereferencing and deproceduring must be applied to a datum of type I in order to yield a datum of type J.

The SAMETYP function also incorporates algorithms for looking up the lists assigned to mode indicants appearing in I or J and for replacing these mode indicants in either structure by pointers to these lists. No attempt is made to find reduced versions of structure declarations. However, when two structures are compared that are structurally equivalent as defined in the ALGOL 68 report, a link to the smaller of the two structures replaces the larger structure in the compiler mode table. This structure comparison and replacement algorithm is essentially the same as the one given by Koster (11) and discussed in Goos (9).

Intermediate Language Generated by the Compiler

Since this report is the beginning of an attempt at producing a formal specification of ALGOL 68 that is an alternative to the Van Wijngaarden and the Vienna notations (12), it was decided that the code produced by our compiler subroutines would be in a systems language rather than in the C.D.C. COMPASS assembly language. The systems language chosen is the PILOT language of M.H. Halstead, versions of which have been implemented on the UNIVAC 1102, the C.D.C. 6000 series of computers, the IBM 360/44, and miscellaneous other machines. PILOT is a

much-distilled subset of the original NELIAC language (13), with some useful features, such as the intermixing of PILOT and machine code where desirable in a program, the use of machine addresses for indirect referencing of variable names, and the addition of partial-word masking operators. In PILOT, commas are used to separate statements, expressions can consist of at most two operands separated by an operator, and assignment of value is to the right, rather than to the left, as in ALGOL 60 or FORTRAN. A typical sequence of PILOT statements might be the following:

```

0 → INDEX REGISTER 1,
U - 1 → U,
LABEL 1: INDEX REGISTER 1 + 1 → INDEX REGISTER 1,
INDEX REGISTER 1 > SIZE : LABEL 2.;;
LHCURR + 1 → [U + INDEX REGISTER 1],
LABEL 1.,
LABEL 2:

```

In the statements above, index register 1 is first initialized to zero, and then a loop is entered at label 1. In the loop, index register 1 is compared to the variable "SIZE". If greater than "SIZE", the statement "LABEL 2." (go to LABEL 2) is executed. Otherwise, execution continues with the value of "LHCURR + 1" stored into the memory location whose address is "U + INDEX REGISTER 1". Further examples of PILOT code are given in Section VII.

Preprocessing Implied by Rules of the Grammar

When a context-free grammar is actually used as the input program of a compiler-writing system, it has to be written with greater care than one ordinarily suspects. In the first place, what is to be done with a grammar

having reserved words that resemble variable names? In Appendix 1, we see that reserved words like "REAL", "PROC", "CASE", etc. could just as well be interpreted as program variables. Next, when we look more closely, we see that the syntax for <name> involves individual letters and digits, whereas reserved words appear as groups of letters. Necessarily, the compiler generated from the grammar above will expect to scan successive words of memory, some of which contain only single letters or digits and some of which contain entire reserved words. Hence, a preprocessor subroutine is needed to read in programs to the compiler.

Along with these trivial operations of packing reserved program words for the compiler, the Appendix 1 syntax calls for the preprocessor to supply a priority digit for each expression operator encountered in a program. Since our grammar assigns unique denotations to operators, the preprocessor need only maintain a table of declared operator priorities in each scanned program block. From this table, a priority digit is inserted following each operator in programs read into the compiler. Naturally, priority declarations are no longer needed by the compiler, and so the preprocessor does not supply them to the compiler.

Stack Mechanisms Used by Compiler Subroutines

As in other compilers for block-structure programming languages, our compiler subroutines use tables to store information about currently valid program variables, labels, data types, and operators. In addition, there are two stacks whose purpose is to facilitate our single-pass translation scheme. The first stack consists of two one-dimensional FORTRAN arrays using the same index, "I". These arrays are referred to as "N(I)" and "TYPE(I)", respectively. As can be seen in the compiler subroutines,

the N-stack is used for constructing the compiler representations of data structures, and the TYPE stack carries auxiliary information concerning whether a given declaration is virtual, actual, or formal, and whether it is stowed or nonstowed. Thus, the combined N-TYPE stack carries context information concerning the syntactic objects recognized by the compiler.

"CØDELØC" is the system name for the second stack, a one-dimensional array with index K. As its name implies, the CØDELØC stack saves information about where nested portions of a program appear in its translation. This information is used selectively by the compiler subroutines for inserting and deleting sections of compiled code in the translated program after the compiler has scanned the corresponding segment of the input program. As will be seen in Section III, it is this use of CØDELØC which enables us to produce reasonably efficient and non-redundant code in a single-pass compilation process. Of course, CØDELØC is also used to supply "target locations" for the jumps implicit in the translation of conditional statements.

Using the N-stack alone, it is easy to demonstrate that a translation grammar specifies a compiler that accepts context-sensitive languages like $\{a^n b^n c^n : n > 0\}$. Next, with the addition of the CØDELØC stack, we see that the resulting compiler system potentially has the same computational ability as a Turing machine. This is because any one of the compiler subroutines could be written so as not to return control to the compiler, and could continue operation by manipulating the two stacks as though they were the tape of a Turing machine. Since the compiler system can be made as general as a Turing machine, any multi-pass compiling algorithm can in theory be written in a single-pass version for a translation grammar. Our goal is to write an efficient and rapid single-pass version of what is usually considered a multi-pass compiler system.

Criteria for Selecting Rule Forms in a Translation Grammar

The criteria for selecting certain rule forms in Appendix 1 rather than others have principally to do with the one-pass nature of the compiler-writing system chosen. For example, a one-pass algorithm cannot efficiently tolerate temporary ambiguities in any subtree of a program. As a consequence, the following rules were chosen that are restrictions of the ALGOL 68 language:

```

<mode indicant> → $<name>$
<operator> → + | - | ... | † <name> †
<call> → <primary> (<actual parameters>)
<slice> → <primary> [<indexers>]
<base> → GØ TØ <label>

```

In the first two rules above, we have forced mode indicants and invented operators to be different in appearance from each other. The next two rules make it impossible to confuse a subscripted variable with a procedure call, and the last rule assures us that the statement

```
v: = if a then b else c fi;
```

does not mean the same thing as

```
v: = if a then b else go to c fi;
```

which is currently a valid interpretation in the ALGOL 68 report.

Again, the desire for an uncomplicated single-scan compiler led us to abandon the ALGOL 68 definition of a block (given in Section 6.1.1 of the report) in favor of a restricted subset of the definition. In the full ALGOL 68 version of a block, the declaration prelude terminates at the first labeled statement or the first statement followed by a jump. To translate such a declaration prelude properly, the compiler needs to know at the beginning of each statement in a declaration prelude what

comes after that statement. Such knowledge could be gained by a separate pass or by requiring forward scans before each statement, but this extra work consumes compiler time that could be better spent doing translation. Hence, our version of the ALGOL 68 block (rules P₃₄ and P₃₅ in Appendix 1) is a subset of the full version in which the declaration prelude terminates at the first executable statement in the block. No generality is lost in this version by not interspersing statements and declarations in the declaration prelude; rather, these interspersed statements are made to appear within the declarations in the blocks that legally may appear there.

The preceding problems of declaration preludes and uniquely identifiable mode and operator indicants are all examples of avoiding syntaxes that result in nondeterministic compilers (6). Another example of this problem is one that arose when an attempt was made to transcribe a portion of the ALGOL 68 report directly into the notation of Appendix 1. In this case the syntax for <stowed declarer> initially took the form:

```
<stowed declarer> → STRUCT (<fields>)
                    | [<rowers>] STRUCT (<fields>)
                    | [<rowers>] <nonstowed declarer>
                    | [<rowers>] <mode indicant>
```

With this set of rules, we are left with three separate methods to declare an array and no way of knowing initially which compiler subroutines should be called. Because of this, the following version of the syntax of arrays and structures appears in Appendix 1:

```

<stowed declarer> → STRUCT (<fields>)
    | <array generator> STRUCT (<fields>)
    | <array generator> <nonstowed declarer>
    | <array generator> <mode indicant>
<array generator> → [<rowers>]

```

Apart from the necessity of avoiding nondeterminisms in the compiler, the other main criterion for selecting certain configurations of rules in a translation grammar is to permit syntax-directed modifications of compiler tables and syntax-directed tests for consistency of program subtrees with non-syntactic objects. As an example of simple modifications made possible by syntax, the following translation rules cause an extra level of reference to be prefixed on the compile-time representation of the declarer in question:

```

P29a: <reference to actual declarer> → <actual mode  

    declarer> / $\psi_{29a1}$ /
P2a : <nonstowed declarer> → REF/ $\psi_{2a1}$ / <mode declarer> / $\psi_{2a2}$ /
P30a: <reference to mode global generator> → HEAP/ $\psi_{30a1}$ /
    <actual mode declarer> / $\psi_{30a2}$ /

```

In the second example, the rules

```

P21a: <virtual mode declarer> → <mode declarer> / $\psi_{21a}$ /
P22a: <actual mode declarer> → <mode declarer> / $\psi_{22a}$ /

```

do not introduce any syntactic ambiguity in the language, since the corresponding transitions of the compiler can only occur in mutually exclusive contexts. However, compiler subroutines ψ_{21a} and ψ_{22a} both test the N-TYPE stack to see that all arrays declared within the <mode declarer> are respectively virtual or actual.

Lastly, we see that syntactic rules may be chosen to make the work of compiler subroutines easier. Good examples of this are the sets of rules below:

$$P_{5a}: \langle \text{fields} \rangle \rightarrow \langle \text{field} \rangle / \psi_{5a1} /$$

$$P_{5b}: \langle \text{fields} \rangle \rightarrow \langle \text{field}, \rangle \langle \text{fields} \rangle / \psi_{5b1} /$$

$$P_{6a}: \langle \text{field}, \rangle \rightarrow \langle \text{field} \rangle, / \psi_{6a} /$$

Subroutine ψ_{6a} generates a compile-time list cell for each field in a structure declaration, and subroutine ψ_{5b1} links together the list cells generated in this fashion. Because of the form of P_{6a} , the structure fields are stored in sequence on the N-TYPE stack by ψ_{6a} . Rule P_{5b} calls for the compiler to work backwards, linking together the two structure fields on top of the N-TYPE stack by successive calls on subroutine ψ_{5b1} . Thus the syntax directs the compiler to link together the compile-time representation of the structure.

On the other hand, the rules

$$P_{7a}: \langle \text{rowers} \rangle \rightarrow \langle \text{rower} \rangle / \psi_{7a1} /$$

$$P_{7b}: \langle \text{rowers} \rangle \rightarrow \langle \text{rowers} \rangle, \langle \text{rower} \rangle / \psi_{7b1}$$

do nothing more than construct a run-time array descriptor in the order of appearance of the program.

III. Translation of Array and Structure Declarations

Preliminaries

ALGOL 68 uses two reservoirs of consecutive memory words for use in generating arrays and structures during program execution. In the first reservoir system, array storage is taken from a stack (called loc in the defining report) whose contents are divided into regions corresponding to currently active program blocks. As the most recent block is abandoned by the program at run time, the storage of its corresponding region in loc is freed for reuse, exactly as in an ALGOL 60 system. As a consequence, structures and arrays stored in loc are only accessible during the lifetime of the block in which they were created. In order to permit a longer lifetime of usage for selected arrays and structures, an alternative reservoir system is available, called the heap in the defining report. The heap system provides that an array or structure is only reclaimed when it can no longer be reached from an active program variable. The use of this heap reservoir necessitates use of a conventional list-processing storage reclamation scheme at run time.

During compilation, the compiler subroutines keep track of whether a currently translated declaration is to be stored in loc or heap. The bookkeeping system used for this purpose is quite simple. The FORTRAN variable "LHCURR" is initialized to contain the characters "LOC", and is temporarily reset to contain the characters "HEAP" whenever compiler subroutines ψ_{30a1} and ψ_{30a2} in translation rule P_{30a} are called. Then, during the translation process, whenever the code for a storage reservation is put out, the contents of LHCURR (denoted in our notation below by "{LHCURR}") is inserted in the translated program.

Array Declarations

There are two parts to the problem of generating arrays in a compiled ALGOL 68 program. The first part involves construction of a run-time descriptor for the array, and the second part involves possible assignment of structures to each element of the resulting array. Translation rules P_{4a} and P_{7a} of Appendix 1 outline the construction of the descriptor, while rules P_{10a} through P_{13e} keep track of whether the descriptor is virtual (has no numerical bounds), actual (has only numerical bounds) or formal (has some mixture of numerical and non-numerical bounds). Since no code is produced for virtual or formal descriptors, compiler subroutine ψ_{4a1} records in the CODELOC stack where translated code for the descriptor begins, so that ψ_{4a2} can erase any code that may have been produced while scanning the <array generator>.

For the recognized program sequence

[<rowers>]

the compiler produces the following translated code:

```

SAVEU,
2 + {LHCURR} → U,
{LHCURR} + 1 + no. of
dimensions → {LHCURR},
0 → [U],
Translation of <rowers>,
U - 1 + no. of
dimensions → TEMP1,
0 → [TEMP1],
[U](36 → 53) → SIZE,
SIZE > 0: U → [TEMP1](0 → 17),
{LHCURR} + SIZE → {LHCURR};;
no. of
dimensions → [TEMP1](36 → 53),
POPUP,

```

The run-time variable U in the above code is used by the compiler sub-routines as an index for placing information in the current word of the array descriptor being generated. As can be seen in Figure 2, the descriptor consists of $n+1$ words, with n the number of dimensions. Thus, the first descriptor word storing bounds is located at address " $2 + \{LHCURR\}$ " above (see Preliminaries for meaning of $\{LHCURR\}$), and that word is set to zero by the " $0 \rightarrow [U]$ " statement.

The statement that reserves space for the descriptor in $\{LHCURR\}$ is inserted later by subroutine ψ_{4a2} after the number of dimensions is known by the compiler. Translation of $\langle \text{rowers} \rangle$ involves calculation of the "strides" of the descriptor. Each stride is associated with a pair of bounds and is the coefficient in a run-time storage-mapping function used in subscripting. When the final stride of the descriptor is calculated, it is stored in the next location following the last pair of bounds in memory and coincides with the actual size of the run-time array. Thus, the code

$$[U](36 + 53) \rightarrow \text{SIZE},$$

above extracts this value and uses it above in a conditional statement that places a link to the first array element into the run-time descriptor:

$$\text{SIZE} > 0: U \rightarrow [\text{TEMP1}](0 \rightarrow 17), \{LHCURR\} + \text{SIZE} \rightarrow \{LHCURR\};;$$

Figure 2 below displays the storage structure used for run-time array descriptors. The " $\ell_i u_i$ " notation represents a two-bit code for indicating whether or not the corresponding lower and upper bounds can be changed dynamically. When ℓ_i (or u_i) is zero, the lower (or upper) bound may vary; otherwise it is fixed. Given this information, we can discuss the translation of $\langle \text{rowers} \rangle$.

Figure 2. Run-Time Representation of Array Descriptors

Type	No. of Dimensions	Block Number	Pointer to First Element of Array
$l_1 u_1$	stride 1	upper bound 1	lower bound 1
.	.	.	.
.	.	.	.
.	.	.	.
$l_n u_n$	stride n	upper bound n	lower bound n

Rules P_{7a} and P_{7b} of Appendix 1 are given below:

$$P_{7a} : \langle \text{rowers} \rangle + \langle \text{rower} \rangle / \psi_{7a1} /$$

$$P_{7b} : \langle \text{rowers} \rangle + \langle \text{rowers} \rangle, \langle \text{rower} \rangle / \psi_{7b1} /$$

Subroutine ψ_{7a1} is called after the translation of the first pair of bounds in an array declaration, and subroutine ψ_{7b1} is called for each subsequent pair of bounds in the declaration. Primarily, ψ_{7a1} and ψ_{7b1} work together to count dimensions and calculate the strides of the descriptor. They assume that translation rules P_{13} , P_{12} , P_{11} , and P_{10} have inserted the upper and lower bounds of each dimension into the appropriate fields of the descriptor word. The code put out by ψ_{7b1} gives an idea of this interaction:


```

[U](18 + 35) - [U](0 + 17) + SIZE,
(Subtracts the lower bound from the upper bound.)
SIZE * [U](36 + 53) + SIZE,
(Calculates the value of the next stride.)
U + 1 + U,
0 + [U],
(Zeroes the next word in the descriptor.)
SIZE > 0 : SIZE + [U](36 + 53);;
(Stores the stride in the next word of the descriptor.)

```

When the array descriptor is completed, there may follow code for initializing each element of the resulting run-time array. This code is written out for rules P_{3b} and P_{3d} :

```

P3b: <stowed declarer> + <array generator>/ $\psi_{3b1}$ /
      STRUCT(/ $\psi_{3b2}$ / <fields>)/ $\psi_{3b3}$ 
P3d: <stowed declarer> + <array generator>/ $\psi_{3d1}$ /
      <mode indicant>/ $\psi_{3d2}$ /

```

Essentially, subroutines ψ_{3d1} and ψ_{3d2} set up a loop that links each element of the generated array to a fresh copy of the <mode indicant>. This copy of the <mode indicant> is created by treating the <mode indicant> as a function whose value is the address of the next location of the storage reservoir specified by LHCURR. Since translation of P_{3b} works in a very similar fashion, we will give as our illustration the code produced by ψ_{3d1} and ψ_{3d2} :

```

0 → INDEX REGISTER 2,
U - 1 → U,
LABEL 2:
INDEX REGISTER 2 + 1 → INDEX REGISTER 2,
INDEX REGISTER 2 > SIZE: LABEL 3.;;
{LHCURR} + 1 → {U + INDEX REGISTER 2},
Translation of <mode indicant>,
(This includes a jump to the definition of the
<mode indicant> and return.)
LABEL 2.,
LABEL 3:

```

Note in the above code that "U - 1" denotes the first address in the region of memory reserved for the stored array, whereas {LHCURR} was set equal to the last address of this region by compiler subroutine ψ_{4a2} . Thus, the first address in any data structure created by the <mode indicant> is given by "{LHCURR} + 1", and this address is the value used to initialize the corresponding array elements iteratively.

Structure Declarations

In our implementation of ALGOL 68, the fields of data structure can store plain values or they can be linked to other structures and arrays. During a declaration, then, code linking a structure field to another structure or array may or may not be called for. On a syntactic basis, we say that such linking code is called for in a structure having some field that corresponds to a <stowed declarer> or <mode indicant>. The code is not compiled in the case of a structure field that corresponds to a <nonstowed declarer>. To simplify our discussion of the translation process, we consider only rule $P_{3,1}$ of Appendix 1, since it uses

essentially the same mechanisms as rule P_{3b} :

$$P_{3a}: \text{<stowed declarer>} \rightarrow \text{STRUCT}(\psi_{3a1}/\text{<fields>}/\psi_{3a2}/$$

The code generated by compiler subroutines ψ_{3a1} and ψ_{3a2} is as follows:

```
SAVE U,
{LHCURR}  $\rightarrow$  U,
{LHCURR} + no. of fields  $\rightarrow$  {LHCURR},
Translation of <fields>,
POP U,
```

The compiler subroutines used above in translating <fields> also count the number of fields in the structure. When the count is completed, subroutine ψ_{3a2} is called, and ψ_{3a2} inserts the field count into the instruction preceding the translation of <fields>. Thus, the code produced by ψ_{3a1} and ψ_{3a2} reserves space for the translated structure in LOC or HEAP and initializes the index of structure fields to point to the first field in the structure to be translated. The code put out for a typical "stowed" field of the structure is produced by compiler subroutines ψ_{5a1} or ψ_{6a} as follows:

```
no. of intervening nonstowed fields + U  $\rightarrow$  U,
{LHCURR} + 1  $\rightarrow$  [U],
Translation of the <field>,
```

In the code above, the index U retains the address of the most recently initialized field in the compiled structure. In order to initialize the next field, U is incremented, and the address of the first free word in LOC or HEAP is inserted in the word whose address is U. Thus, the run-time field is made to point to the first address in the generated structure corresponding to the translated <field>. Of course, no code at all will appear for a "nonstowed" field.

IV. Translation of Mode Declarations

As seen in rule P_{25a} , a mode declaration has the form

$M\text{ODE } \langle \text{mode indicant} \rangle = \langle \text{actual mode declarer} \rangle.$

Sample code produced for such a declaration by compiler subroutines

ψ_{25a1} , ψ_{25a2} , and ψ_{25a3} is as follows:

```

    LABEL 5.,
    LABEL 4:
    Translation of  $\langle \text{actual mode declarer} \rangle,$ 
    [RETURN] .,
    LABEL 5:

```

In the above code, "LABEL 4" is the unique program label corresponding to the definition of the $\langle \text{mode indicant} \rangle$. "LABEL 5" is used to isolate the indicant definition in the run-time program. With this isolation, the only way to execute the code produced for $\langle \text{actual mode declarer} \rangle$ is to jump to LABEL 4. After this declarer is executed at run-time, there follows the instruction

```
[RETURN] .,
```

which is used in the PILOT language for a jump to the location whose address is stored in RETURN.

In order to execute this data-definitional function, we need a function call in the form of a $\langle \text{mode indicant} \rangle$ appearing in some declaration.

This function call produces the following code:

```

    SAVE RETURN,
    J + 2 + RETURN,
    LABEL 4.,
    POP RETURN

```

For this calling sequence, LABEL 4 is used as the location of the particular <mode indicant> to be called. The PILOT run-time program counter stored in J is saved in the RETURN variable so that the code for the mode definition can jump back to the statement "POP RETURN". All the compiler needs to do is carry a record of the labels that correspond to the currently active mode indicants and to retain a representation of the data structures to which these indicants correspond. As before, the value of this mode indicant function is the first free location in LOC or HEAP at the calling point in the compiled program.

A Final Example

To bring together the information given in our previous descriptions, we present a typical mode declaration, together with its generated code:

The Declaration

```
MODE $A$ = STRUCT([1:5] STRUCT($B$ X, REAL Y)Z);
```

The Compiled Code

```

LABEL 10.,
LABEL 7:
SAVE U,
LOC → U,
LOC + 1 → LOC,
U + 1 → U,
LOC → [U],
SAVE U,
2 + LOC → U,
LOC + 2 → LOC,
0 → [U],
1 → [U](35 + 53),
1 → OPERAND REGISTER 2,
1 → [U](55 + 55),
OPERAND REGISTER 2 → [U](0 + 17),
5 → OPERAND REGISTER 2,
1 → [U](54 + 55),
OPERAND REGISTER 2 → [U](18 + 35),
1 → [U](36 + 53),
[U](18 + 35) - [U](0 + 17) → SIZE,
SIZE + 1 → SIZE,
U + 1 → U,
0 → [U],
SIZE > 0: SIZE → [U](36 + 53);;
```

```

U - 2 + TEMP1,
O + [TEMP1],
[U](36 + 53) + SIZE,
SIZE > 0: U + [TEMP1](0 + 17),
LOC + SIZE + LOC;;
1 + [TEMP1](36 + 53),
POP U,
O + INDEX REGISTER 3,
U-1 + U,
LABEL 8:
INDEX REGISTER 3 + 1 + INDEX REGISTER 3,
INDEX REGISTER 3 > SIZE: LABEL 9.;;
LOC + 1 + [U + INDEX REGISTER 3],
SAVE U,
LOC + U,
LOC + 2 + LOC,
1 + U + U,
LOC + 1 + [U],
SAVE RETURN,
J + 2 + RETURN,
LABEL {B$}.,
POP RETURN
POP U,
LABEL 8.,
LABEL 9:
POP U,
[RETURN] .,
LABEL 10:

```

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Appendix 1

A Partial Translation Grammar for ALGOL 68*

$P_{1a} : \langle \text{mode declarer} \rangle \rightarrow \langle \text{stowed declarer} \rangle / \psi_{1a} /$
 $P_{1b} : \quad \quad \quad | \langle \text{nonstowed declarer} \rangle / \psi_{1b} /$
 $P_{1c} : \quad \quad \quad | \langle \text{mode indicator} \rangle / \psi_{1c} /$
 $P_{2a} : \langle \text{nonstowed declarer} \rangle \rightarrow \text{REF} / \psi_{2a1} / \text{mode declarer} / \psi_{2a2} /$
 $P_{2b} : \quad \quad \quad | \text{mode declarator} / \psi_{2b} /$
 $P_{3a} : \langle \text{stowed declarer} \rangle \rightarrow \text{STRUCT} (/ \psi_{361} / \langle \text{fields} \rangle) / \psi_{3a2} /$
 $P_{3b} : \quad \quad \quad | \langle \text{arraygenerator} \rangle / \psi_{361} / \text{STRUCT} (/ \psi_{362} / \langle \text{fields} \rangle) / \psi_{363} /$
 $P_{3c} : \quad \quad \quad | \langle \text{arraygenerator} \rangle * \text{nowstowed declarer} / \psi_{3c} /$
 $P_{3d} : \quad \quad \quad | \langle \text{arraygenerator} \rangle / \psi_{3d1} / \langle \text{mode indicator} \rangle / \psi_{3d2} /$
 $P_{4a} : \langle \text{arraygenerator} \rangle \rightarrow [/ \psi_{4a1} / \text{rowers}] / \psi_{4a2} /$
 $P_{5a} : \langle \text{fields} \rangle \rightarrow \langle \text{field} \rangle / \psi_{5a1} /$
 $P_{5b} : \quad \quad \quad | \langle \text{field} \rangle , / \psi_{5b1} / \langle \text{fields} \rangle / \psi_{5b2} /$
 $P_{6a} : \langle \text{field} \rangle \rightarrow \langle \text{field mode declarer} \rangle \langle \text{field selector} \rangle / \psi_{6a} /$
 $P_{7a} : \langle \text{rowers} \rangle \rightarrow \langle \text{rower} \rangle / \psi_{7a1} /$
 $P_{7b} : \quad \quad \quad | \langle \text{rowers} \rangle , \langle \text{rower} \rangle / \psi_{7b1} /$
 $P_{8a} : \langle \text{field mode declarer} \rangle \rightarrow \langle \text{mode declarer} \rangle$
 $P_{9a} : \langle \text{field selector} \rangle \rightarrow \langle \text{name} \rangle / \psi_{9a} /$
 $P_{10a} : \langle \text{rower} \rangle \rightarrow : / \psi_{10a} /$
 $P_{10b} : \quad \quad \quad | \langle \text{lower bound} \rangle : \langle \text{upper bound} \rangle$
 $P_{11a} : \langle \text{lower bound} \rangle \rightarrow \langle \text{bound} \rangle / \psi_{11a} /$
 $P_{12a} : \langle \text{upper bound} \rangle \rightarrow \rangle \langle \text{bound} \rangle / \psi_{12a} /$
 $P_{13a} : \langle \text{bound} \rangle \rightarrow \langle \text{formula} \rangle / \psi_{13a} /$
 $P_{13b} : \quad \quad \quad | \text{EITHER} / \psi_{13b} /$
 $P_{13c} : \quad \quad \quad | \text{FLEX} / \psi_{13c} /$

* This syntax has been tested, and is certified to be LR(1).

$P_{13}:$ |<formula EITHER/ Ψ_{13d} /

$P_{13e}:$ | formula FLEX/ Ψ_{13e} /

$P_{14a}:$ <mode decarator> \rightarrow <procdecl>

$P_{14b}:$ |<union mode decl>

$P_{14c}:$ | REAL/ Ψ_{14c} /

$P_{14d}:$ | INT/ Ψ_{14d} /

$P_{14e}:$ | B00L/ Ψ_{14f} /

$P_{14f}:$ | CHAR/ Ψ_{14f} /

$P_{15a}:$ <procdecl> \rightarrow PROC/ Ψ_{15a} /

$P_{15b}:$ | PROC/ Ψ_{15b1} /<virtual procplan>/ Ψ_{15b2} /

$P_{16a}:$ <virtual procplan> \rightarrow (<virtual parameters>)/ Ψ_{16a} /

$P_{16b}:$ |<virtual mode declarer>/ Ψ_{16b} /

$P_{16c}:$ |(<virtual parameters>)<virtual mode declarer>/ Ψ_{16c} /

$P_{17a}:$ <virtual parameters> \rightarrow <virtual mode declarer>/ Ψ_{17a} /

$P_{17b}:$ |<virtual mode declarer>, Ψ_{17b1} /<virtual parameters>/ Ψ_{17b2} /

$P_{18a}:$ <union mode decl> \rightarrow UNION (/ Ψ_{18a1} /<nonunion indicants>)/ Ψ_{18a2} /

$P_{19a}:$ <nonunion indicants> \rightarrow <mode indicant>/ Ψ_{19a} /

$P_{19b}:$ |<mode indicant>, Ψ_{19b1} /<nonunion indicants>/ Ψ_{19b2} /

$P_{20a}:$ <mode indicant> \rightarrow \$ <name> \$/ Ψ_{20a} /

$P_{21a}:$ <virtual mode declarer> \rightarrow <mode declarer>/ Ψ_{21a} /

$P_{22a}:$ <actual mode declarer> \rightarrow mode declarer/ Ψ_{22a} /

$P_{23}:$ <formal mode declarer> \rightarrow <mode declarer>

$P_{24a}:$ <declaration> \rightarrow <mode declaration>

$P_{24b}:$ |<operator declaration>

$P_{24c}:$ |<initialization>

$P_{24d}:$ |<priority declaration>

$P_{25a}:$ $\langle \text{mode declaration} \rangle \rightarrow \text{MODE} / \psi_{25a1} / \langle \text{mode indicant} \rangle = / \psi_{25a2} /$
 $\quad \quad \quad \langle \text{actual mode declarer} \rangle / \psi_{25a3} /$

$P_{26a}:$ $\langle \text{operator declaration} \rangle \rightarrow \text{OP} / \psi_{26a1} / \langle \text{op.} \rangle = / \psi_{25a2} / \langle \text{routine} \rangle / \psi_{26a3} /$

$P_{27a}:$ $\langle \text{initialization} \rangle \rightarrow \langle \text{reference to actual declarer} \rangle / \psi_{27a} / \langle \text{name list} \rangle$

$P_{27b}:$ $\quad \quad \quad | \langle \text{reference to mode global generator} \rangle / \psi_{27b} / \langle \text{name list} \rangle$

$P_{27c}:$ $\quad \quad \quad | \langle \text{formal mode declarer} \rangle \langle \text{tag} \rangle = / \psi_{27c1} / \langle \text{tertiary} \rangle / \psi_{27c2} /$

$P_{28a}:$ $\langle \text{name list} \rangle \rightarrow \langle \text{tag} \rangle / \psi_{28a} /$

$P_{28b}:$ $\quad \quad \quad | \langle \text{tag} \rangle : = / \psi_{28b1} / \langle \text{tertiary} \rangle / \psi_{28b2} /$

$P_{28c}:$ $\quad \quad \quad | \langle \text{name list} \rangle, \langle \text{tag} \rangle / \psi_{28c} /$

$P_{28d}:$ $\quad \quad \quad | \langle \text{name list} \rangle, \langle \text{tag} \rangle : = / \psi_{28d1} / \langle \text{tertiary} \rangle / \psi_{28d2} /$

$P_{29a}:$ $\langle \text{reference to actual declarer} \rangle \rightarrow \langle \text{actual mode declarer} \rangle / \psi_{29a} /$

$P_{30a}:$ $\langle \text{reference to mode global generator} \rangle \rightarrow \text{HEAP} / \psi_{30a1} / \langle \text{actual mode declarer} \rangle /$
 $\quad \quad \quad \psi_{30a2} /$

$P_{31a}:$ $\langle \text{program} \rangle \rightarrow \$\text{ENTRY}\$ / \psi_{31a1} / \langle \text{particular program} \rangle \$\text{EXIT}\$ / \psi_{31a2} /$

$P_{32a}:$ $\langle \text{particular program} \rangle \rightarrow \langle \text{closed clause} \rangle$

$P_{33a}:$ $\langle \text{closed clause} \rangle \rightarrow \text{BEGIN} / \psi_{33a1} / \langle \text{serial clause} \rangle \text{END} / \psi_{33a2} /$

$P_{33b}:$ $\quad \quad \quad | (/ \psi_{33a1} / \langle \text{serial clause} \rangle) / \psi_{33a2} /$

$P_{33c}:$ $\quad \quad \quad | (/ \psi_{33a1} / \langle \text{row of clause} \rangle) / \psi_{33a2} /$

$P_{34a}:$ $\langle \text{serial clause} \rangle \rightarrow \langle \text{declaration prelude sequence} \rangle ; \langle \text{suite of clause train} \rangle$

$P_{34b}:$ $\quad \quad \quad | \langle \text{declaration prelude sequence} \rangle ; \langle \text{label completer} \rangle \langle \text{suite of clause train} \rangle$

$P_{34c}:$ $\quad \quad \quad | \langle \text{label completer} \rangle \langle \text{suite of clause train} \rangle$

$P_{34d}:$ $\quad \quad \quad | \langle \text{suite of clause train} \rangle$

$P_{35a}:$ $\langle \text{declaration prelude} \rangle \rightarrow \langle \text{unitary clause} \rangle ; \langle \text{declaration prelude} \rangle$

$P_{35b}:$ $\quad \quad \quad | \langle \text{declaration} \rangle$

$P_{36a}:$ $\langle \text{declaration prelude sequence} \rangle \rightarrow \langle \text{declaration prelude sequence} \rangle ;$

$P_{36b}:$ $\langle \text{unitary clause} \rangle \rightarrow \langle \text{unitary clause} \rangle ;$

$P_{36c}:$ $\langle \text{unitary clause EXIT} \rangle \rightarrow \langle \text{unitary clause} \rangle \text{EXIT}$

$P_{36d}:$ $\langle \text{label} : \rangle \rightarrow \langle \text{label} \rangle :$

$P_{37a} :$ $\langle \text{suite of clause train} \rangle \rightarrow \langle \text{unitary clause} \rangle .$
 $P_{37b} :$ $| \langle \text{labelstat} \rangle$
 $P_{37c} :$ $| \langle \text{labelstat} \rangle \langle \text{suite of clause train} \rangle$
 $P_{37d} :$ $| \langle \text{unitary clause}; \rangle \langle \text{suite of clause train} \rangle$
 $P_{38a} :$ $\langle \text{labelstat} \rangle \rightarrow \langle \text{unitary clause}; \rangle \langle \text{label} : \rangle$
 $P_{38b} :$ $| \langle \text{unitary clause EXIT} \rangle \langle \text{label} : \rangle$
 $P_{38c} :$ $| \langle \text{labelstat} \rangle \langle \text{label} : \rangle$
 $P_{38d} :$ $\langle \text{label completer} \rangle \rightarrow \langle \text{label} ? \rangle$
 $P_{38e} :$ $| \langle \text{label completer} \rangle \langle \text{label} : \rangle$
 $P_{38f} :$ $\langle \text{declaration prelude sequence} \rangle \rightarrow \langle \text{declaration prelude} \rangle / \psi_{34b2} /$
 $P_{38g} :$ $| \langle \text{declaration prelude sequence}; \rangle \langle \text{declaration prelude} \rangle$
 $P_{39a} :$ $\langle \text{label} \rangle \rightarrow \langle \text{name} \rangle$
 $P_{40a} :$ $\langle \text{unitary clause} \rangle \rightarrow \langle \text{tertiary} \rangle$
 $P_{40b} :$ $| \langle \text{confrontation} \rangle$
 $P_{40c} :$ $| \text{FØR} \langle \text{tag} \rangle \text{FRØM} \langle \text{serial clause} \rangle \text{BY} \langle \text{serial clause} \rangle \text{TØ} \langle \text{serial clause} \rangle$
 $\quad \quad \text{DØ} \langle \text{unitary clause} \rangle$
 $P_{40d} :$ $| \text{FRØM} \langle \text{serial clause} \rangle \text{BY} \langle \text{serial clause} \rangle \text{TØ} \langle \text{serial clause} \rangle \text{DØ}$
 $\quad \quad \langle \text{unitary clause} \rangle$
 $P_{40e} :$ $| \text{WHILE} \langle \text{serial clause} \rangle \text{DØ} \langle \text{unitary clause} \rangle$
 $P_{41a} :$ $\langle \text{confrontation} \rangle \rightarrow \langle \text{identity relation} \rangle$
 $P_{41b} :$ $| \langle \text{conformity relation} \rangle$
 $P_{41c} :$ $| \langle \text{reference to mode assignation} \rangle$
 $P_{41d} :$ $| \langle \text{mode cast} \rangle$
 $P_{42a} :$ $\langle \text{reference to mode assignation} \rangle \rightarrow \langle \text{reference to mode tertiary} \rangle$
 $\quad \quad : = \langle \text{unitary clause} \rangle$
 $P_{43a} :$ $\langle \text{reference to mode tertiary} \rangle \rightarrow \langle \text{tertiary} \rangle$
 $P_{44a} :$ $\langle \text{conformity relation} \rangle \rightarrow \langle \text{reference to mode tertiary} \rangle$
 $\quad \quad \langle \text{conformity relator} \rangle \langle \text{tertiary} \rangle$

P_{45a}: <conformity relator> → ::
 P_{45b}: | ::=
 P_{46a}: <identity relations> → <reference to mode tertiary>
 <identity relator><reference to mode tertiary>
 P_{47a}: <identity relator> → ::=
 P_{47b}: | ::=:
 P_{48a}: <mode cast> → <virtual mode declarer>/_{48a1}/ <unitary clause>/_{48a2}/
 P_{49a}: <void cast> → :<unitary clause>/_{49a}/
 P_{50a}: <tertiary> → <formula>
 P_{51a}: <formula> → <formula><op.><unary>
 P_{51b}: | <unary>
 P_{52a}: <unary> → <op.><unary>
 P_{52b}: | <secondary>
 P_{53a}: <priority declaration> → PRIORITY<op>=<nonzero digit>
 P_{54a}: <moid cast> → <mode cast>
 P_{54b}: | <void cast>
 (Rules P₅₅ through P₆₀ of first version are deleted.)
 P_{61a}: <op.> → +
 P_{61b}: | -
 P_{61c}: | *
 P_{61d}: | /
 P_{61e}: | / ./.
 P_{61f}: | +
 P_{61g}: | ^
 P_{61h}: |
 P_{61i}: | \lceil
 P_{61j}: | =

$P_{61k}:$ | >
 $P_{61l}:$ | <
 $P_{61m}:$ | \geq
 $P_{61n}:$ | \leq
 $P_{61o}:$ | =
 $P_{61p}:$ | \neq
 $P_{61q}:$ | -=
 $P_{61r}:$ | +=
 $P_{61s}:$ | *=
 $P_{61t}:$ | /=
 $P_{61u}:$ | ./.=
 $P_{61v}:$ | RE
 $P_{61w}:$ | IM
 $P_{61x}:$ | ØDD
 $P_{61y}:$ | SIGN
 $P_{61z}:$ | RØUND
 $P_{61aa}:$ | \uparrow <name> \uparrow
 $P_{62a}:$ <secondary> \rightarrow <primary>
 $P_{62b}:$ {<cohesion>
 $P_{63a}:$ <cohesion> \rightarrow <selection>
 $P_{63b}:$ }<generator>
 $P_{64a}:$ <selection> \rightarrow <name> ØF <secondary>
 $P_{65a}:$ <generator> \rightarrow <reference to mode global generator>
 $P_{65b}:$ |<reference to mode local generator>
 $P_{66a}:$ |<reference to mode local generator> \rightarrow LOC <actual mode declarer>
 $P_{67a}:$ <primary> \rightarrow <base>
 $P_{67b}:$ |<closed clause>
 $P_{67c}:$ |<conditional clause>

$P_{68a}:$ $\langle \text{base} \rangle \rightarrow \langle \text{tag} \rangle$
 $P_{68b}:$ $|\langle \text{denotation} \rangle$
 $P_{68c}:$ $|\langle \text{slice} \rangle$
 $P_{68d}:$ $|\langle \text{call} \rangle$
 $P_{68e}:$ $|\text{GØ TØ } \langle \text{label} \rangle$
 $P_{68f}:$ $|\text{SKIP}$
 $P_{68g}:$ $|\text{NIL}$
 $P_{69a}:$ $\langle \text{tag} \rangle \rightarrow \langle \text{name} \rangle$
 $P_{70a}:$ $\langle \text{name} \rangle \rightarrow \langle \text{letter} \rangle$
 $P_{70b}:$ $|\langle \text{name} \rangle \langle \text{letter} \rangle$
 $P_{70c}:$ $|\langle \text{name} \rangle \langle \text{digit} \rangle$
 $P_{70d}:$ $|\langle \text{name} \rangle *$
 $P_{71a}:$ $\langle \text{denotation} \rangle \rightarrow \langle \text{number} \rangle$
 $P_{71b}:$ $|\langle \text{character denotation} \rangle$
 $P_{71c}:$ $|\langle \text{string denotation} \rangle$
 $P_{71d}:$ $|\langle \text{bits} \rangle$
 $P_{71e}:$ $|\langle \text{routine} \rangle$
 $P_{72a}:$ $\langle \text{routine} \rangle \rightarrow (*(\langle \text{formal parameters} \rangle) \langle \text{mold case} \rangle *)$
 $P_{72b}:$ $|\langle \text{mold cast pack} \rangle$
 $P_{73a}:$ $\langle \text{mold cast pack} \rangle \rightarrow (* / \Psi_{73a1} / \langle \text{mold cast} \rangle *) / \Psi_{73a2} /$
 $P_{74a}:$ $\langle \text{formal parameters} \rangle \rightarrow \langle \text{formal parameters} \rangle$
 $P_{74b}:$ $|\langle \text{formal parameter}; \rangle \quad \langle \langle \text{formal parameter} \rangle;$
 $P_{75a}:$ $\langle \text{formal parameter}; \rangle \rightarrow \langle \text{formal parameter} \rangle;$
 $P_{76a}:$ $\langle \text{formal parameter} \rangle \rightarrow \langle \text{formal mode declarer} \rangle \langle \text{name} \rangle$
 $P_{77a}:$ $\langle \text{call} \rangle \rightarrow \langle \text{primary} \rangle (/ \Psi_{77a1} / \langle \text{actual parameters} \rangle) / \Psi_{77a2} /$
 $P_{78a}:$ $\langle \text{actual parameters} \rangle \rightarrow \langle \text{unitary clause} \rangle / \Psi_{78a1} /$
 $P_{78b}:$ $|\langle \text{unitary clause} \rangle, / \Psi_{78b1} / \langle \text{actual parameters} \rangle / \Psi_{78b2} /$
 $P_{79a}:$ $\langle \text{slice} \rangle \rightarrow \langle \text{primary} \rangle [\langle \text{indexer} \rangle]$

$P_{80a}:$ $\langle \text{indexer} \rangle \rightarrow \langle \text{trimscript} \rangle$
 $P_{80b}:$ $|\langle \text{indexer} \rangle, \langle \text{trimscript} \rangle$
 $P_{81a}:$ $\langle \text{trimscript} \rangle \rightarrow \langle \text{sum} \rangle$
 $P_{81b}:$ $|\text{ : AT } \langle \text{sum} \rangle$
 $P_{81c}:$ $|\langle \text{sum} \rangle \text{ : } \langle \text{sum} \rangle \text{ AT } \langle \text{sum} \rangle$
 $P_{81d}:$ $|\langle \text{sum} \rangle \text{ : AT } \langle \text{sum} \rangle$
 $P_{81e}:$ $|\langle \text{sum} \rangle \text{ : } \langle \text{sum} \rangle$
 $P_{81f}:$ $|\text{ : } \langle \text{sum} \rangle \text{ AT } \langle \text{sum} \rangle$
 $P_{82a}:$ $\langle \text{conditional clause} \rangle \rightarrow \text{IF}/\psi_{33a1}/\langle \text{serial clause} \rangle$
 $\text{THEN}/\psi_{82a2}/\langle \text{serial clause} \rangle$
 $\text{ELSE}/\psi_{82a3}/\langle \text{serial clause} \rangle \text{FI}/\psi_{82a4}/$
 $P_{82b}:$ $|\text{IF}/\psi_{33a1}/\langle \text{serial clause} \rangle \text{THEN}/\psi_{82a2}/\langle \text{serial clause} \rangle \text{FI}/\psi_{82b3}/$
 $P_{82c}:$ $|\text{CASE}/\psi_{33a1}/\langle \text{serial clause} \rangle \text{IN}/\psi_{82d2}/\langle \text{row of clause} \rangle \text{ESAC}/\psi_{82c3}/$
 $P_{82d}:$ $|\text{CASE}/\psi_{33a1}/\langle \text{serial clause} \rangle \text{IN}/\psi_{82d2}/\langle \text{row of clause} \rangle \text{ELSE}/\psi_{82d3}/$
 $\langle \text{serial clause} \rangle \text{ESAC}/\psi_{82d4}/$
 $P_{83a}:$ $\langle \text{row of clause} \rangle \rightarrow \langle \text{unitary clause}, \rangle \langle \text{unitary clause} \rangle / \psi_{83a}/$
 $P_{83b}:$ $|\langle \text{unitary clause}, \rangle \langle \text{row of clause} \rangle / \psi_{83b}/$
 $P_{84a}:$ $\langle \text{unitary clause}, \rangle \rightarrow \langle \text{unitary clause}, \rangle / \psi_{84a}/$
 $P_{85a}:$ Deleted
 $P_{86a}:$ $\langle \text{number} \rangle \rightarrow \langle \text{a integer} \rangle$
 $P_{86b}:$ $|\langle \text{a integer} \rangle . \langle \text{b integer} \rangle$
 $P_{86c}:$ $|\langle \text{a integer} \rangle .$
 $P_{86d}:$ $|\langle \text{b integer} \rangle$
 $P_{87a}:$ $\langle \text{a integer} \rangle \rightarrow \langle \text{nonzero digit} \rangle$
 $P_{87b}:$ $|\langle \text{a integer} \rangle \langle \text{digit} \rangle$
 $P_{88a}:$ $\langle \text{b integer} \rangle \rightarrow \langle \text{digit} \rangle$
 $P_{88b}:$ $|\langle \text{b integer} \rangle \langle \text{digit} \rangle$
 $P_{89a}:$ $\langle \text{bits} \rangle \rightarrow 0 \langle \text{octal digit} \rangle$

P_{90a}: <octal integer> → <octal digit>
P_{90b}: | <octal integer> <octal digit>
P_{91a}: <octal digit> → 0
P_{91b}: | <nonzero octal>
P_{92a}: <nonzero octal> → 1
.
.
.
.
P_{92g}: | 7
P_{93a}: <digit> → 0
P_{93b}: | <nonzero digit>

SUBROUTINE CLEAN	10
ENTRY PSI1A	20
;	30
ENTRY PSI1C	40
;	50
TYPE(I)=6RSTOWED	60
RETURN	70
;	80
ENTRY PSI1B	90
;	100
TYPE(I)=9RNONSTOWED	110
RETURN	120
;	130
ENTRY PSI2A1	140
;	150
K=K+1	160
CODELOC(K)=Z	170
RETURN	180
;	190
ENTRY PSI2A2	200
C	210
C CODE FOR THIS NONSTOWED DECLARER IS DESTROYED AT THIS POINT. INFORMATI	220
C ABOUT FORMAL DECLARERS IS PASSED ON.	230
C	240
IF (LINK(CONT(N(I)),1).EQ.3RFRM) GO TO 101	250
ITEMP1=3RVCT	260
GO TO 102	270
101 ITEMP1=3RFRM	280
102 ITEMP=IFETCH(1)	290
CALL STRIND (4LOREF,ITEMP)	300
CALL SETIND (N(I),ITEMP,2)	310
CALL SETIND (ITEMP1,ITEMP,1)	320
N(I)=ITEMP	330
Z=CODELOC(K)	340
CALL REDUCE (K)	350
RETURN	360
C	370
ENTRY PSI2B	380
C	390
CALL SETIND (3RVCT,N(I),1)	400
RETURN	410
C	420
ENTRY PSI3A1	430
C	440
ENTRY PSI3B2	450
C	460
C	470
C A STORAGE RESERVATION COMMAND FOR THE STRUCTURE IS INSERTED HERE	480
C BY SUBROUTINE PSI3A2	490
C	500
K=K+1	510
CODELOC(K)=Z+3	520
K=K+1	530
CODELOC(K)=Z+4	540
FIELDNO(K)=1	550


```

CODELOC(K)=CONVERT(S,2R L) 1110
CALL OUTCDA2 (Z,CODELOC(K),10H; 1120
CALL OUTCDA4 (Z,ITEMP,5H + 1,ITEMP,10H, 1130
S=S+1 1140
CALL OUTCDA4 (Z,ITEMP,9H > SIZE 1,CODELOC(K-1),10H. 1150
CALL OUTCDA4 (Z,LHCURR,10H+1 + (U + ,ITEMP,10H), 1160
RETURN 1170
ENTRY PSI3B3 1180
1190
1200
THISROUTINE IS ALMOST PSI3A2 FOLLOWED BY PSI3D2,WITHFACILITIES FOR 1210
CARRYING OVER INFORMATION ABOUT VIRTUAL ORACTUALDECLARATIONS. 1220
1230
ITEMP=IFETCH(1) 1240
CALL STRIND (4LOSTR,ITEMP) 1250
CALL SETIND (N(I),ITEMP,2) 1260
CALL SETINC (LINK(CONT(N(I))+1),1),ITEMP,1) 1270
COUNTNO=IBYTE(CONT(N(I)),10) 1280
CALL PUTIND (COUNTNO,ITEMP,10) 1290
N(I)=ITEMP 1300
CALL REDUCE (K) 1310
INDEX=CODELOC(K) 1320
CALL OUTCDB6 (INDEX,LHCURR,5H +,CONVERT(COUNTNO,2R ),5H +,L 1330
LHCURR,10H, 1340
IF (COUNTNO.GT.FIELDNO(K)) GO TO 105 1350
INDEX=CODELOC(K)-2 1360
CALL OUTCDB1 (INDEX,2H//) 1370
CALL OUTCDP1 (INDEX,2H//) 1380
GO TO 106 1390
105 CALL OUTCDA1 (Z,10HPOPU, 1400
106 CALL REDUCE (K) 1410
1420
CALL OUTCDA2 (Z,CODELOC(K),10H. , 1430
K=K-1 1440
CALL OUTCDA2 (Z,CODELOC(K),10H; 1450
CALL REDUCE (K) 1460
CALL SETIND (N(I),N(I-1),2) 1470
1480
AT THIS POINT, INFORMATION ABOUT WHETHER THE STRUCTURE IS VIRTUAL IS 1490
COMBINED WITH THE SAME INFORMATION CONCERNING THE ARRAY GENERATOR. 1500
1510
ITEMP=LINK(CONT(N(I)),1) 1520
ITEMP1=LINK(CONT(N(I-1)),1) 1530
IF (ITEMP.EQ.3RACT.AND.ITEMP1.EQ.3RACT) GO TO 108 1540
IF (ITEMP.EQ.3RFRM.OR.ITEMP1.EQ.3RFRM) GO TO 107 1550
CALL SETIND (3RVRT,N(I-1),1) 1560
GO TO 108 1570
107 CALL SETIND (3RFRM,N(I-1),1) 1580
108 I=I-1 1590
R=R-1 1600
RETURN 1610
1620
ENTRY PSI3C 1630
1640
AT THIS POINT, ARRANGEMENTS ARE MADE FOR SUBSCRIPTING OF BIT AND 1650

```

```

C CHARACTER ARRAYS. THERE IS NO NEED TO ERASE ANY CODE FOR NONSTOWED 1660
C DECLARERS, SINCE RULES PSI2A2, PSI14, PSI18, AND PSI15 DO THIS. 1670
C 1680
      ITEMPL=LINK(CONT(N(I),3)) 1690
      IF (ITEMP.EQ.3RBIT) GO TO 110 1700
      IF (ITEMP.EQ.3RBYT) GO TO 112 1710
109 CALL SETIND (N(I),N(I-1),2) 1720
      I=I-1 1730
      RETURN 1740
110 INC=10H/60 * (U - 1750
111 R=R+1 1760
      S=S+1 1770
      IF (R.GT.MAXR) MAXR=R 1780
      ITEMPL=CONVERT(R,2RIR) 1790
      ITEMPL1=CONVERT(S,2R L) 1800
      CALL OUTCDA3 (Z,3H0*,ITEMP,10H, 1810
      CALL OUTCDA8 (Z,ITEMPL1,6H* (U- ,ITEMP,9H)(36*53) ,INC,ITEMP,10H)(3 1820
      16*53) , ,10H 1830
      CALL OUTCDA4 (Z,ITEMP,6H +1 *,ITEMP,10H, 1840
      CALL OUTCDA5 (Z,ITEMP,10H<DIMENSION,1H*,ITEMPL1,10H.*; 1850
      R=R-1 1860
      GO TO 109 1870
112 INC=10H12*(U - 1880
      GO TO 111 1890
C 1900
      ENTRY PSI3D2 1910
C 1920
      CALL OUTCDA2 (Z,CODELOC(K),10H. , 1930
      K=K-1 1940
      CALL OUTCDA (Z,CODELOC(K),10H* 1950
      CALL REDUCE (K) 1960
      CALL SETIND (N(I),N(I-1),2) 1970
      I=I-1 1980
      R=R-1 1990
      RETURN 2000
C 2010
      ENTRY PSI4A1 2020
C 2030
      K=K+1 2040
      CODELOC(K)=Z 2050
      FIELDNO(K)=0 2060
      CALL OUTCDA1 (Z,10HSAVEU, 2070
      CALL OUTCDA2 (Z,LHCURR,10H+ 2 * U 2080
      Z=Z+1 2090
      CALL OUTCDA2 (Z,9H00 * (U),,10H 2100
C 2110
C THIS OUTPUT CODE MAKES U THE INDEX OF THE DESCRIPTOR BEING CREATED. 2120
C A SPACE IS LEFT IN THE OUTPUT CODE FOR LATER INSERTION OF A STATEMENT 2130
C OF THE FORM LHCURR+DIMENSION* LHCURR THAT RESERVES SPACE FOR THE 2140
C ARRAY DESCRIPTOR ON EITHER THE HEAP OR THE LOG STACK. RULE PSI4A2 2150
C MAKES THIS INSERTION 2160
C 2170
      RETURN 2200
C 2210
      ENTRY PSI4A2 2220

```

```

C
C CAUSES THE TOTAL SIZE OF THE ARRAY TO BE CALCULATED, SEES WHETHER THE
C ARRAY IS VIRTUAL, AND STORES THE DIMENSION IN THE DESCRIPTOR.
C
    BOUNDNO=LINK(FIELDNO(K),1)
    I=I+1
    N(I)=IFETCH(1)
    CALL PUTIND (BOUNDNO,N(I),10)
    CALL SETIND (3RARY,N(I),3)
    VIRTUAL=LINK(FIELDNO(K),2)
    IF (VIRTUAL.EQ.0) GO TO 115
    IF (VIRTUAL.EQ.BOUNDNO) GO TO 113
    CALL SETIND (3RFRM,N(I),1)
C
C THE ARRAY IS FORMAL.
C
    GO TO 114
113 CALL SETIND (3RVRT,N(I),1)
C
C THE ARRAY IS VIRTUAL, AND ITS ASSOCIATED TRANSLATION IS DESTROYED.
C
114 Z=CODELOC(K)
    K=K-1
    RETURN
115 CALL SETIND (3RACT,N(I),1)
C
C FOR AN ACTUAL ARRAY, WE PROCEED TO PUT OUT CODE FOR COMPLETING
C THE DESCRIPTOR.
C
    ITEM=CODELOC(K)+3
    ITEM1=CONVERT(BOUNDNO,2R )
C
C WE NOW HAVE THE INFORMATION THAT PSI7B1 STORED THE SIZE OF THE ARRAY
C IN U AT RUN TIME, WHERE U IS CURRENTLY THE ADDRESS OF THE WORD THAT
C FOLLOWS THE DESCRIPTOR IN MEMORY.
C
    CALL OUTCDB6 (ITEM,LHCURR,4H +,ITEM1,4H *,LHCURR,5H, //)
    CALL REDUCE (K)
    CALL OUTCDA4 (Z,3HU -,ITEM1,8H *,TEMP1,10H, //)
    CALL OUTCDA2 (Z,10H00*[TEMP1],10H, //)
    CALL OUTCDA2 (Z,10H(U)(36*53),10H*SIZE, //)
    CALL OUTCDA3 (Z,7HSIZE>0?,10HU*[TEMP1](,10H0*17), //)
    CALL OUTCDA4 (Z,LHCURR,10H + SIZE *,LHCURR,10H;: //)
    CALL OUTCDA3 (Z,ITEM1,10H-1*[TEMP1],10H(36*53),//)
C
C AT THIS POINT, U IS THE ADDRESS OF THE FIRST ARRAY ELEMENT.
C
    RETURN
C
    ENTRY PSI5A1
C
    ENTRY PSI6A
C
C CHECKS THE STRUCTURE FIELD TO DISCOVER WHETHER STOWED OR NONSTOWED. IF
C NONSTOWED, DELETES PREVIOUSLY GENERATED CODE FOR POINTING FROM FIELD

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```

: TO AVAILABLE STORAGE AT RUN TIME, AND INCREMENTS A GENERAL COUNT OF      2780
: FIELDS. IF STOWED OR MODE INDICANT, PUTS OUT CODE FOR INCREMENTING RUN    2790
: TIME ALLOCATION POINTER TO GENERATE A CONNECTION BETWEEN A FIELD AND      2800
: ITS DATA STRUCTURE.                                                     2810
:                                                                           2820
:   IF (TYPE(I).EQ.6RSTOWED) GO TO 116                                     2830
:   Z=CODELOC(K)-1                                                         2840
:   FIELDNO(K)=1+FIELDNO(K)                                               2850
:   GO TO 117                                                             2860
116  ITEMP=CODELOC(K)                                                     2870
:   CALL OUTCDP2 (ITEMP,CONVERT(FIELDNO(K),2R ),10H+U+U,  //)           2880
:   CALL OUTCDP2 (ITEMP,LHCURR,10H+U,  //)                               2890
:   FIELDNO(K)=1                                                         2900
117  CODELOC(K)=Z+1                                                       2910
:   Z=Z+2                                                                2920
:   RETURN                                                                2930
:                                                                           2940
:   ENTRY PSI5B1                                                         2950
:                                                                           2960
: LINKS TOGETHER TWO FIELDS OF A STRUCTURE AT COMPILE TIME. COUNTS THE    2970
: NUMBER OF FIELDS. COMBINES INFORMATION ABOUT ACTUAL AND VIRTUAL FIELDS  2980
:                                                                           2990
:   CALL SETIND (N(I),N(I-1),2)                                          3000
:   CALL PUTIND (1+IBYTE(CONT(N(I))),10),N(I-1),10)                     3010
:   ITEMP=LINK(CONT(1+N(I)),1)                                           3020
:   ITEMP1=LINK(CONT(N(I-1)),1)                                           3030
:   IF (ITEMP.EQ.3RACT.AND.ITEMP1.EQ.3RACT) GO TO 120                    3040
:   IF (ITEMP.EQ.3RFRM.OR.ITEMP1.EQ.3RFRM) GO TO 118                    3050
:   IF (ITEMP.EQ.3RVCT.OR.ITEMP1.EQ.3RVCT) GO TO 119                    3060
:   CALL SETIND (3RVCT,N(I-1),1)                                          3070
:   GO TO 120                                                             3080
118  CALL SETIND (3RFRM,N(I-1),1)                                          3090
:   GO TO 120                                                             3100
119  CALL SETIND (3RVCT,N(I-1),1)                                          3110
120  I=I-1                                                                3120
:   RETURN                                                                3130
:                                                                           3140
:   ENTRY PSI7A1                                                         3150
:                                                                           3160
:   CALL SETDIR (1,FIELDNO(K),1)                                          3170
:   CALL OUTCDA2 (Z,10H01+U(36+10H53),  //)                             3180
:   CALL OUTCDA4 (Z,10H(U)(18+35),10H-[U](0+17),6H+SIZE,,10HU+1+U,  //  3190
:   1)                                                                    3200
:   CALL OUTCDA4 (Z,5H00+U,,7HSIZE+0+10HSIZE+U(3,10H6+53) ;; //)       3210
:   RETURN                                                                3220
:                                                                           3230
:   ENTRY PSI7B1                                                         3240
:                                                                           3250
: COUNTS UP THE NUMBER OF DIMENSIONS, STORES THE STRIDE IN THE CURRENT    3260
: DESCRIPTOR WORD, AND ZEROES THE NEXT DESCRIPTOR WORD.                  3270
:                                                                           3280
:   CALL SETDIR (1+LINK(FIELDNO(K),1),FIELDNO(K),1)                    3290
:   CALL OUTCDA3 (Z,10H(U)(18+35),10H-[U](0+17),10H+SIZE,  //)         3300
:                                                                           3310
: CALCULATES UPPERBOUND-LOWERBOUND.                                       3320

```

CALL OUTCDA3 (Z,10HSIZE*(U1(3,10H6*53)*SIZE,10H, //)	3330
CALL OUTCDA1 (Z,10HU*1*U, //)	3340
7=Z+1	3350
CALL OUTCDA4 (Z,6H0*(U),,7HSIZE>0,10HSIZE*(U1(3,10H6*53): //)	3360
RETURN	3370
	3380
ENTRY PSI8A	3390
	3400
GENERATES A CELL IN THE COMPILE-TIME LIST COPY OF A STRUCTURE	3410
DECLARATION.	3420
	3430
ITEMP=IFETCH(2)	3440
CALL STRIND (N(I),ITEMP+1)	3450
I=I-1	3460
IF (TYPE(I).EQ.6RSTOWED) GO TO 121	3470
CALL STRIND (4L0+FL,ITEMP)	3480
GO TO 122	3490
121 CALL STRIND (4L1+FL,ITEMP)	3500
122 CALL SETINC (N(I),ITEMP,2)	3510
N(I)=ITEMP	3520
RETURN	3530
	3540
ENTRY PSI9A	3550
	3560
I=I+1	3570
N(I)=NAME	3580
RETURN	3590
	3600
ENTRY PSI10A	3610
	3620
: THIS ROUTINE CCUNTS VIRTUAL BOUNDS IN AN ARRAY DECLARATION.	3630
: CALL SETDIR (1+LINK(FIELDNO(K),2),FIELDNO(K),2)	3640
: NOTE THAT ANYTHING BUT ZERO IN THE SECOND FIELD OF FIELDNO INDICATES	3650
: THAT THE ARRAY IS NOT AN ACTUAL ARRAY.	3660
: RETURN	3670
: ENTRY PSI11A	3680
: ITEMP1=10H(U1(55*55)	3690
: ITEMP2=10H(U1(0*17)	3700
: GO TO 123	3710
: ENTRY PSI12A	3720
: ITEMP1=10H(U1(54*54)	3730
: ITEMP2=10H(U1(18*35)	3740
123 GO TO (124,126,126,125,125), BOUND	3750
124 CALL OUTCDA3 (Z,5H 01* ,ITEMP1,10H, //)	3760
125 CALL OUTCDA4 (Z,CONVERT(P,2R T),1H*,ITEMP2,10H, //)	3770
P=P-1	3780
I=I-1	3790
	3800
	3810
	3820
	3830
	3840
	3850
	3860
	3870

126 RETURN	3880
ENTRY PSI13B	3890
BOUND=2	3900
RETURN	3910
ENTRY PSI13C	3920
BOUND=3	3930
RETURN	3940
ENTRY PSI13A	3950
BOUND=1	3960
GO TO 127	3970
ENTRY PSI13D	3980
BOUND=4	3990
GO TO 127	4000
ENTRY PSI13E	4010
BOUND=5	4020
127 IF (SAMETYP(N(I),INTEGER,REF)) GO TO 128	4030
CALL ERROR(110HNON-INTEGER AFRAY BOUND IN BOUNDLIST.	4040
1 RETURN	4050
128 CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z)	4060
RETURN	4070
ENTRY PSI14C	4080
ITEMP=3RREA	4090
GO TO 130	4100
ENTRY PSI14D	4110
ITEMP=3RINT	4120
GO TO 130	4130
ENTRY PSI14E	4150
ITEMP=3RRIT	4160
GO TO 130	4170
ENTRY PSI14F	4180
ITEMP=3RBYT	4190
GO TO 130	4200
ENTRY PSI15A	4210
129 ITEMP=3RPRC	4220
	4230
	4240
	4250
	4260
	4270
	4280
	4290
	4300
	4310
	4320
	4330
	4340
	4350
	4360
	4370
	4380
	4390
	4400
	4410

130 I=I+1	4420
ITEMP2=IFETCH(1)	4430
CALL SETIND (ITEMP,ITEMP2,3)	4440
RETURN	4450
ENTRY PSI15B1	4460
K=K+1	4470
CODELOC(K)=7	4480
GO TO 129	4490
ENTRY PSI15B2	4500
CODE GENERATED DURING ANALYSIS OF THE <VIRTUAL PROCPLAN> IS DESTROYED.	4510
Z=CODELOC(K)	4520
CALL REDUCE (K)	4530
RETURN	4540
ENTRY PSI16A	4550
CALL PUTIND (1+IBYTE (CONT(N(I)),10),N(I),10)	4560
ITEMP1=2	4570
GO TO 131	4580
ENTRY PSI16B	4590
ITEMP1=1	4600
131 ITEMP=IFETCH(1)	4610
CALL SETIND (N(I),ITEMP,ITEMP1)	4620
I=I-1	4630
CALL SETIND (ITEMP,N(I),2)	4640
RETURN	4650
ENTRY PSI16C	4660
ITEMP=IFETCH(1)	4670
CALL SETIND (N(I),ITEMP,1)	4680
CALL SETIND (N(I-1),ITEMP,2)	4690
CALL PUTIND (IBYTE (CONT(N(I-1)),10),N(I),10)	4700
I=I-2	4710
N(I)=ITEMP	4720
RETURN	4730
ENTRY PSI17A	4740
ITEMP1=1	4750
GO TO 132	4760
ENTRY PSI17B1	4770
ITEMP1=0	4780
132 ITEMP=IFETCH(1)	4790
CALL SETIND (3RPRM,ITEMP,3)	4800
CALL SETIND (N(I),ITEMP,2)	4810
	4820
	4830
	4840
	4850
	4860
	4870
	4880
	4890
	4900
	4910
	4920
	4930
	4940
	4950
	4960

CALL PUTIND (ITEMP1,ITEMP,10)	4970
N(I)=ITEMP	4980
RETURN	4990
ENTRY PSI1782	5000
THIS ROUTINE LINKS TOGETHER THE PARAMETERS IN THE COMPILER COPY OF THE	5010
PARAMETER LIST,AND COUNTS THE NUMBER OF PARAMETERS.	5020
CALL PUTIND (1+IBYTE(CONT(N(I)),10),N(I-1),10)	5030
CALL SETIND (N(I),N(I-1),1)	5040
I=I-1	5050
RETURN	5060
ENTRY PSI18A1	5070
K=K+1	5080
CODELOC(K)=Z	5090
RETURN	5100
ENTRY PSI18A2	5110
Z=CODELOC(K)	5120
CALL REDUCE (K)	5130
RETURN	5140
ENTRY PSI19A	5150
ENTRY PSI19B1	5160
ITEMP=IFETCH(1)	5170
CALL SETIND (3RUNI,ITEMP,3)	5180
CALL SETIND (N(I),ITEMP,2)	5190
N(I)=ITEMP	5200
RETURN	5210
ENTRY PSI19B2	5220
CALL SETIND (N(I),N(I-1),1)	5230
I=I-1	5240
RETURN	5250
ENTRY PSI20A	5260
I=I+1	5270
ITEMP=IFETCH(2)	5280
N(I)=ITEMP	5290
CALL STRIND (NAME,ITEMP+1)	5300
CALL SETIND (3RMOD,ITEMP,3)	5310
NAME=ALIAS(NAME)	5320
THE ALIAS FUNCTION LOOKS UP THE RUN-TIME PROGRAM LABEL ASSIGNED TO THE	5330
MODE INDICANT,AND ADDS THE INDICANT TO THE COMPILER DECLARATION TABLE	5340
IF NECESSARY.	5350
	5360
	5370
	5380
	5390
	5400
	5410
	5420
	5430
	5440
	5450
	5460
	5470
	5480
	5490
	5500
	5510

```

      CALL OUTCDA4(Z,10HSAVE RETUR,2HN,,10HJ+2*RETURN,3H,/)
      CALL OUTCDA4 (Z,NAME,2H.,,10HPOP RETURN,10H,      //)
      RETURN
C
      ENTRY PSI21A
C
C A CHECK IS MADE HERE TO SEE WHETHER THE DECLARER IN QUESTION CAN BE
C USED AS A VIRTUAL DECLARER. NO CHECK IS NEEDED FOR FORMAL DECLARERS.
C
      ITEM=LINK(CONT(N(I)),1)
      IF (ITEM.EQ.3RVRT.OR.ITEM.EQ.3RVCT) RETURN
      CALL ERROR(110HAN ATTEMPT WAS MADE TO USE A FORMAL OR ACTUAL D
1DECLARER IN PLACE OF A VIRTUAL DECLARER. )
      RETURN
C
      ENTRY PSI22A
C
C A CHECK IS MADE TO SEE WHETHER THE DECLARER IN QUESTION CAN BE USED AS
C AN ACTUAL DECLARER. NO CHECK IS NEEDED FOR FORMAL DECLARERS.
C
      ITEM=LINK(CONT(N(I),1))
      IF (ITEM.EQ.3RACT.OR.ITEM.EQ.3RVCT) RETURN
      CALL ERROR(110HAN ATTEMPT WAS MADE TO USE A FORMAL OR VIRTUAL
1DECLARER IN PLACE OF AN ACTUAL DECLARER. )
      RETURN
C
      ENTRY PSI25A1
C
      ENTRY PSI26A1
C
      K=K+1
      CODELOC(K)=Z
C
C THIS RULE ALLOWS DESTRUCTION OF CODE PRODUCED BY RULE PSI20A.
C
      RETURN
C
      ENTRY PSI25A2
C
      Z=CODELOC(K)+2
C
C A SPACE IS RESERVED FOR A TRANSFER OF CONTROL AROUND THE MODE
C DECLARATION IN THE TRANSLATED PROGRAM.
C
      CALL OUTCDB2 (Z,NAME,10H:      //)
C
C THE MODE INDICANT WAS DECLARED BY PSI20A, AND NAME CONTAINS THE RUN-
C TIME LABEL ASSOCIATED WITH THE INDICANT.
C
      FIELDNO(K)=MODE
C
C MODE IS THE INDEX OF THE COMPILER MODES TABLE.
C
      RETURN
C

```

5530
 5540
 5550
 5560
 5570
 5580
 5590
 5600
 5610
 5620
 5640
 5650
 5660
 5670
 5680
 5690
 5700
 5710
 5720
 5740
 5750
 5760
 5770
 5780
 5790
 5800
 5810
 5820
 5830
 5840
 5850
 5860
 5870
 5880
 5890
 5900
 5910
 5920
 5930
 5940
 5950
 5960
 5970
 5980
 5990
 6000
 6010
 6020
 6030
 6040

```

ENTRY PSI25A3
CALL OUTCDA3 (Z,9HRETURN+RJ,8H(30+47),,10HRJ.,    //)
S=S+1
ITEMP1=CONVERT(S,2R L)
ITEMP=CODELOC(K)+1
CALL OUTCOR2 (ITEMP,ITEMP1,10H.,    //)
CALL OUTCDA2 (Z,ITEMP1,10H:    //)
ITEMP=FIELDNO(K)
MODES(ITEMP,3)=N(I)
ASSIGNS COMPILER DATA STRUCTURE TO MODE INDICANT.
I=I-1
CALL REDUCE (K)
RETURN
ENTRY PSI26A2
AT THIS POINT, THE OPERATOR HAS BEEN ENTERED INTO THE OPS TABLE. THE
CURRENT INDEX OF OPS IS OP, AND THE RUN-TIME ALIAS OF THE OPERATOR IS
STORED IN THE COMPILER VARIABLE NAME, EXACTLY AS FOR THE CASE OF MODE
INDICANTS.
Z=CODELOC(K)+2
A SPACE IS RESERVED FOR A TRANSFER OF CONTROL AROUND THE OPERATOR
DEFINITION IN THE TRANSLATED PROGRAM.
CALL OUTCOR2 (Z,NAME,10H:    //)
FIELDNO(K)=OP
RETURN
ENTRY PSI26A3
CALL OUTCDA3 (Z,9HRETURN+RJ,8H(30+47),,10HRJ.,    //)
S=S+1
ITEMP1=CONVERT(S,2R L)
ITEMP=CODELOC(K)+1
CALL OUTCOR2 (ITEMP,ITEMP1,10H.,    //)
ITEMP=FIELDNO(K)
OPS(ITEMP,3)=N(I)
CALL OUTCDA2 (Z,ITEMP1,10H:    //)
I=I-1
CALL REDUCE (K)
RETURN
C
ENTRY PSI27A
C THIS ROUTINE ADDS AN EXTRA LEVEL OF REFERENCING TO THE DECLARER, BOTH I
C COMPILER REPRESENTATION AND THE RUN-TIME CODE9 LHCURR CARRIES INFORMA
C ABOUT WHETHER LOC OR HEAP IS IN USE.
S=S+1
ITEMP=CONVERT(S,2R L)

```

```

6050
6060
6070
6080
6090
6100
6110
6120
6130
6140
6150
6160
6170
6200
6210
6220
6230
6240
6250
6260
6270
6280
6290
6300
6310
6320
6330
6340
6350
6360
6370
6380
6390
6400
6410
6420
6430
6440
6450
6460
6470
6480
6490
6500
6510
6520
6530
6540
6550
6560
6570
6580
6590
6600
6610

```

```

FIELDNO(K)=ITEMP                                6620
ITEMP1=CODELOC(K)+1                             6630
CALL OUTCDP6 (ITEMP1,ITEMP,4H,1+,LHCURR,1H+,LHCURR,8H,    //) 6640
CALL OUTCDP4 (ITEMP1,LHCURR,5H+1+,LHCURR,10H),           //) 6650
ITEMP1=ITEMP1-2                                   6660
S=S+1                                              6670
ITEMP=CONVERT(S,2R L)                             6680
133 CALL OUTCDP2 (ITEMP1,ITEMP,10H,                    //) 6690
CALL OUTCDA3 (Z,10HRETURN,RJ,10H(30+47),10HRJ.,        //) 6700
CALL OUTCDA2 (Z,ITEMP,10H,                            //) 6710
ITEMP=IFETCH(1)                                    6720
CALL SETIND (3RREF,ITEMP,3)                        6730
CALL SETIND (N(I),ITEMP,2)                         6740
N(I)=ITEMP                                          6750
RETURN                                              6760
                                                6770
ENTRY PSI27B                                       6780
                                                6790
: LIKE PSI27A,EXCEPT THAT THE EXTRA LEVEL OF RUN-TIME INDIRECTNESS IS 6800
: SUPPLIED BY PSI30A1.                            6810
:                                                  6820
: S=S+1                                           6830
: ITEMP=CONVERT(S,2R L)                          6840
: FIELDNO(K)=ITEMP                               6850
: ITEMP1=CODELOC(K)+1                            6860
: CALL OUTCDP2 (ITEMP1,ITEMP,10H,                //) 6870
: S=S+1                                           6880
: ITEMP1=ITEMP1-1                                6890
: ITEMP=CONVERT(S,2R L)                          6900
: GO TO 133                                       6910
:                                                  6920
:                                                  6950
: ENTRY PSI27C1                                  6960
:                                                  6970
: THE TAG RULE PSI68A HAS THE EFFECT OF ENTERING THE NAME ONTO THE NAMES 6980
: TABLE, AND LEAVES THE RUN-TIME ALIAS OF <TAG> IN THE VARIABLE NAME. 6990
: CODE FOR THE VIRTUAL DECLARER IS DESTROYED      7000
:                                                  7010
: OPERAND(I)=NAME                                7020
:   TYPE(I)=9RNONSTORED
: Z=CODELOC(K)                                   7040
: RETURN                                          7050
:                                                  7060
: ENTRY PSI27C2                                  7070
:                                                  7080
: IF (SAMETYP(N(I-1),N(I),REF).AND.REF.EQ.0) GO TO 134 7090
:   CALL ERROR(110)ATTEMPT TO INITIALIZE A DECLARED VARIABLE
:   USING A <TERTIARY> OF A DIFFERENT MODE.      )
:   RETURN                                       7110
134 IF (TYPE(I).NE.9RNONSTORED)GO TO 135          7130
:   P=P+1                                       7140
:   ITEMP=CONVERT(P,2R T)
:   CALL OUTCDA4(Z,OPERAND(I),1H+,ITEMP,10H,      //)
:   OPERAND(I)=ITEMP                            7200
135 I=I-1

```

```

      OPERAND(I)=IHASHB(OPERAND(I),N(I))
;
; SETS THE TYPE OF THE NAME AND THE REGION WHERE ITS VALUE IS STORED.
;
      CALL OUTCDA4(Z,OPERAND(I+1),1H,OPERAND(I),10H,      //)
      RETURN
;
;
      ENTRY PSI28A
      ENTRY PSI28C
;
; ENTER THE<TAG> AND ITS MODE ONTO THE NAME TABLE.
;
      NAME=IHASHB(NAME,N(I))
      CALL OUTCDA4 (Z,LHCURR,4H+1,NAME,10H);      //)
      CALL OUTCDA3 (Z,10HSAVE RETUR,10HN,J+2,RETU,10HRN,      //)
      CALL OUTCDA2 (Z,FIELDNO(K),10H,      //)
      CALL OUTCDA2 (Z,10HPOP RETURN,10H,      //)
;
      ENTRY PSI28B1
      ENTRY PSI28D1
;
      OPERAND(I)=IHASHB(NAME,N(I))
      TYPE(I)=9RNONSTORED
      RETURN
;
      ENTRY PSI28B2
      ENTRY PSI28D2
;
; THE DATA TYPES OF <TERTIARY> AND <TAG> ARE COMPARED. THE <TERTIARY> I
; DEREFERENCED AND DEPROCEDURED TO ONE LEVEL OF REFERENCE BELOW N(I-1)
; NOTE THAT THE FIRST PARAMETER OF SAMETYP IS THE ONE TO BE LOWERED.
;
      IF (SAMETYP(N(I),LINK(CONT(N(I-1)),1),REF)) GO TO 169
      CALL FRROR (110HTHE <TERTIARY> ASSIGNED TO THE <TAG> OF THIS MODE
1DECLARATION DOES NOT POSSESS THE SAME MODE AS ITS ASSIGNEE. )
      RETURN
169 IF (TYPE(I).NE.9RNONSTORED) GO TO 170
      P=P+1
      ITEMP=CONVERT(P,2R T)
      CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP,10H,      //)
170 CALL OUTEVAL (P,REF,OUTCODE,7,1000-Z)
      I=I-1
      CALL OUTCDA4 (Z,ITEMP,1H,OPERAND(I),10H,      //)
      RETURN
;
;
      ENTRY PSI29A1
;
      ITEMP=IFETCH(1)
      CALL SETIND (3RREF,ITEMP,3)
      CALL SETIND (N(I),ITEMP,2)
      N(I)=ITEMP
      RETURN
;
      ENTRY PSI30A1

```

```

. LHCURR=4HHEAP
HERE, WE PUT OUR CODE FOR GENERATING AN EXTRA LEVEL OF REFERENCING AT
RUN TIME. THE BEGINNING IS SAVED FOR USE AS DESTINATION IN ASSIGNMENTS
CALL OUTCDA5 (Z,10HHEAP+1*HEA,10HP,HEAP+1 *,10H[HEAP], 8HSAVEU,,
10HHEAP*U, //)
RETURN
ENTRY PSI30A2
LHCURR=3HLCC
CALL OUTCDA3 (Z,10H U*GENERAT,10HOR VALUE, ,10HPOP U, //)
THE VALUE OF THE GENERATOR IS PRESERVED FOR USE IN ASSIGNMENTS9
OUTCODE(Z,3)=10HPOP U, //
RETURN
ENTRY PSI31A1
CALL INITIAL(NAME,MODE)
RETURN
ENTRY PSI31A2
CALL FINAL(NAME,MODE)
RETURN
RETURN
ENTRY PSI33A1
ENTRY PSI33B1
LHCURR=3HLCC
CALL BLOK3GN(P,MODE)
ITEMP1=Z+3
K=K+1
CODELOC(K)=ITEMP1
K=K+1
CODELOC(K)=ITEMP1
RESERVES SPACE FOR THE TRANSLATION OF INITIALIZATION DECLARATIONS AND
FOR MANAGING RUN-TIME STORAGE ALLOCATION.
Z=Z+4
RETURN
ENTRY PSI33A2
ENTRY PSI33B2
CALL BLOKEND(P,MODE)
RETURN

```

```

7840
7850
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7870
7880
7890
7900
7910
7920
7930
7940
7950
7960
7970
7980
7990
8000
8010
8020
8030
8040
8050
8060
8080
8090
8100
8110
8120
8130
8140
8150
8210
8220
8230
8240
8250
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8270
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8290
8300
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8320
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8340
8350
8360
8430
8440

```

PSI34B CAUSES LOCAL BLOCK STORAGE TO BE ERASED IN THE TRANSLATED PROGRAM.

ENTRY PSI34A

FILLS IN BLANK SPACES AT THE BEGINNING OF TRANSLATED CODE FOR THE CURRENT BLOCK.

```

      ITEMPC=CODELOC(K)-2
      ITEMPC1=ITEMPC+3
      DO 138 ITEMPC2=ITEMPC,ITEMPC1,1
138 OUTCODE(ITEMPC2,1)=2H//
      K=K-1
      RETURN

```

ENTRY PSI34B1

```

      K=K-1
      RETURN

```

ENTRY PSI34B2

CODE IS PUT OUT AT THIS POINT TO FILL IN THE STORAGE MANIPULATION BLANKS LEFT BY PSI33A1 AND PSI33B1.

```

      ITEMPC=CODELOC(K)-2
      CALL OUTCDP2 (ITEMPC,10H1+9BLOKNUM,10HBLOKNUM,/)
      CALL OUTCDP3 (ITEMPC,4HLOC,10HSTORAGE[BL,10HOKNUM],/)
      CALL OUTCDA6 (Z,10HSTORAGE[BL,10HOKNUM]LOC,1H,10HBLOKNUM-1,10HBL
10KNUM, ,10H      /)
      RETURN

```

ENTRY PSI36A

CLEARS THE MOST RECENT DECLARATION FROM THE N STACK AND RESETS CODELOC FOR THE NEXT DECLARATION.

```

      LHCURR=3HLOC
      CODELOC(K)=Z+1
      Z=Z+2
      I=I-1
      RETURN

```

CASE/PSI33A1/<SER. CL.>IN/PSI82D2/<ROW OF CL.>FI/PSI82C3/
CASE/PSI33A1/<SER. CL.>IN/PSI82D2/<ROW OF CL.>ELSE/PSI82D3/<SER. CL.>FI

```

      ENTRY PSI82D2
      IF(.NOT.SAMETYP(N(I),INTREAL,REF))
1      CALL ERROR(110HINDEX OF A CASE STATEMENT IS NOT A NUMERICAL QU
2ANTITY.
      IF (.NOT.(TYPE(I).EQ.9RNONSTORED)) GO TO 183
      P=P+1
      ITEMPC=CONVEPT(P,2R T)

```

8450
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8530
8540
8550
8560
8570
8580
8590
8600
8610
8620
8630
8640
8650
8660
8670
8680
8690
8700
8710
8720
8730
8740
8750
8760
8770
8780
8790
8800
8810
8820
8830
8840
8850
8860
8870
8880
8890
8900
8910
8920
8930

8950
8960
8970


```

CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP,10H,      //)      8980
CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z)          8990
I=I-1                                           9000
CODELOC(K)=Z+1                                 9010
CALL OUTCDA4 (Z,ITEMP,6H  <0,0,10H.;;        //)      9020
CALL OUTCDA6 (Z,ITEMP,6H  > ,0,6H      ,0,10H.;; //)      9030
CALL OUTCDA6 (Z,ITEMP,6H  + ,0,6H      ,ITEMP,10H, //)      9040
CALL OUTCDA3 (Z,ITEMP,10H+RJ(30+47),10H, RJ., //)      9050
                                           9060
: T(P)<0:LABEL ELSE.;;                          9070
: T(P)> NUMBER OF CLAUSES IN ROW: LABEL ELSE.;;  9080
: T(P)+ LABEL FI + T(P), (SEE PSI8203 FOR COMPLETION.) 9090
: T(P)+RJ(30+47),RJ.,                          9100
: CAUSES A JUMP TO A TABLE LEADING TO THE APPROPRIATE CLAUSE OF THE 9110
: CASE STATEMENT, NOTE THAT STORAGE(BLOKNUM) IS USED TO RESET P TO THE 9120
: SAME LEVEL FOR EACH CLAUSE.                  9130
: THE SAVING OF CODELOC(K) IN WHAT FOLLOWS IS PERFORMED FOR COMPATIBILI 9140
: WITH PSI33A1 WHEN A<ROW OF CL.> APPEARS WITHIN A <COLLATERAL CLAUSE>. 9150
                                           9160

      CALL BLOKEND(P,MODE)
      K=K+1          $CODELOC(K)=Z
      Z=Z+1          9220
      K=K+1          9230
      CODELOC(K)=100 9240
      RETURN         9250
:                                     9260
:                                     9270
      ENTRY PSI82D3 . 9280
:                                     9290
: SEARCH HERE FOR THE MOST DEREFERENCED AND DEPROCEDURED (I.E.,LOWEST) 9300
: CLAUSE IN THE <ROW OF CLAUSE>, AND BRING THE REMAINING CLAUSES DOWN 9310
: TO THAT LEVEL.                          9320
                                           9330
      CALL BLOKEND(P,MODE)
      CALL BLOKRGN(P,MODE)
      K=K-2          9350
      S=S-1          9360
      Z=Z-1          9370
      LINK1=MINIMUM(N(I)) 9380
                                           9390
: RULES PSI83A AND PSI83B FOR <ROW OF CLAUSE> CREATE A LIST OF LABELS 9400
: ON TOP OF THE OPERAND STACK.EACH ELT. OF THE LIST CONTAINS IN SEQUENC 9410
: (1) A POINTER TO THE NEXT ELT.          9420
: (2) THE LABEL WHERE THE UNITARY CLAUSE BEGINS. 9430
: (3) THE LABEL WHERE DEREFERENCING AND DEPROCEDURING 9440
: OF THAT CLAUSE OCCURS.                  9450
: (4) THE LOCATION IN THE CODE OF (3) WHERE ESCAPE IS 9460
: MADE TO THE CODE DIRECTLY FOLLOWING THE CASE 9470
: STATEMENT, (SEE BALANCE SUBROUTINE.) 9480
                                           9490
: ITEMp1=OPERAND(I) 9500
: ITEMp2=N(I)       9510
: CALL BALANCE (LINK1,ITEMp1,ITEMp2,S,P,Z,OUTCODE) 9520
: N(I)=LINK1        9530
:                                     9540

```

```

LINK1 POINTS TO THE LOWEST DATA TYPE IN THE ROW OF CLAUSE.
SEE PSI82D2 FOR CONTENTS OF CODELOC(K) AND CODELOC(K-1) HERE.

S=S+1
ITEMP=CONVERT(S,2R L)
CALL OUTCDA2 (Z,ITEMP,10H,      //)
ITEMP1=CODELOC(K)
K=K-1
OUTCODE(ITEMP1+2,3)=ITEMP

THE LOCATION OF THE CASE STATEMENT JUMP TABLE IS PUT INTO THE JUMP COD

ITEMP2=OPERAND(I)
139 CALL OUTCDA2 (Z,CONT(ITEMP2+2),10H,      //)
ITEMP2=LINK(CONT(ITEMP2),1)
IF (ITEMP2.GT.0) GO TO 139

THE JUMP TABLE IS GENERATED. THE ELSE CLAUSE LABEL IS NEXT INSERTED.

S=S+1
ITEMP=CONVERT(S,2R L)
CALL OUTCDA2 (Z,ITEMP,10H,      //)
OUTCODE(ITEMP1,3)=ITEMP
OUTCODE(ITEMP1+1,3)=CONVERT(IRYTE (CONT(OPERAND(I)),10),2R )
OUTCODE(ITEMP1+1,5)=ITEMP
GO TO 184

ENTRY PSI82D4

FURTHER BALANCING TAKES PLACE HERE BETWEEN THE <ROW OF CLAUSE> AND TH
<SERIAL CLAUSE>.
AFTER SAVING THE BLOCK VALUE, DO THE BLOCK END BOOKKEEPING.

CALL BLOKEND(P,MODE)
P=P+1
ITEMP4=CONVERT(P,2R T)
IF (OPERAND(I).EQ.ITEMP4) GO TO 140
CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP4,10H,      //)
140 K=K-1

TEST HERE FOR WHETHER THE <SER. CL.> IS LOWER THAN THE <ROW OF CL.>.
IF LOWER, JUMP AT RUN TIME TO THE CODE FOLLOWING THE CASE STATEMENT,
CAUSE THE <ROW OF CL.> RESULTING TO BE FURTHER DEREFERENCED AND DEPRO
IF HIGHER, DEREFERENCE AND DEPROCEDURE THE <SER. CL.>; THEN INSERT TH
FOR JUMPING OUT OF THE <ROW OF CL.> ROUTINES.

IF (.NOT.SAMETYP(N(I-1),N(I),REF)) GO TO 141
IF (REF.EQ.0) GO TO 142
GO TO 144
141 IF (SAMETYP(N(I),N(I-1),REF)) GO TO 145

CALL ERROR(110)THE REMAINING CLAUSE AFTER*ELSE* IN THE CASE ST
ATEMENT DOES NOT HAVE THE SAME MODE AS PRECEDING CLAUSES.

RETURN

```

```

9550
9560
9570
9580
9590
9600
9610
9620
9630
9640
9650
9660
9670
9680
9690
9700
9710
9720
9730
9740
9750
9760
9770
9780
9790
9800
9810
9820
9830
9840
9850
9930
9960
9890
9900
9910
10000
10010
10020
10030
10040
10050
10060
10070
10080
10090
10100
10110
10130
10140

```

```

NO BALANCING IS NEEDED HERE.
142 S=S+1
    ITEMP=CONVERT(S,2R L)
    CALL OUTCDA2 (Z,ITEMP,10H:      //)
    CALL LOOP82D (OPERAND(I-1),ITEMP,OUTCODE)

    PROVIDES ESCAPE LABELS TO THE CASE CLAUSES.

143 I=I-1
    TYPE(I)=6RSTORED
    OPERAND(I)=ITEMP4
    RETURN

    THE ENTIRE <ROW OF CL.>, AFTER BEING BALANCED INTERNALLY, IS HERE BALAN
    THE SERIAL ELSE CLAUSE.

144 S=S+1
    ITEMP=CONVERT(S,2R L)
    CALL OUTCDA2 (Z,ITEMP,10H:,      //)
    TYPE(I)=ITEMP
    S=S+1
    ITEMP=CONVERT(S,2R L)

    DO LOOP82D TO FORCE ALL <ROW OF CL.> TO ESCAPE TO THIS POINT IN THE
    CODE. NOTE THAT VALUE RETURN IS USED TO CARRY THE VALUES OF EACH CLAU

    CALL LOOP82D (CPERAND(I-1),ITEMP,OUTCODE)

    CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z)
    CALL OUTCDA2 (Z,TYPE(I),10H:      //)
    N(I-1)=N(I)
    GO TO 143

    THE <SERIAL CL.> MUST BE DEREFERENCED AND DEPROCEURED TO BALANCE WITH
    <ROW OF CL.>.

145 S=S+1
    ITEMP=CONVERT(S,2R L)
    CALL OUTCDA2 (Z,ITEMP,10H:,      //)
    CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z)
    CALL LOOP82D (OPERAND(I-1),ITEMP,OUTCODE)
    CALL OUTCDA2 (Z,ITEMP,10H:      //)
    GO TO 143

    ENTRY PSI82C3
    CALL BLOKEND(P,MODE)
    K=K-2
    LINK1=MINIMUM(N(I))
    S=S-1
    Z=Z-1
    CALL BALANCE (LINK1,OPERAND(I),N(I),S,P,Z,OUTCODE)
    N(I)=LINK1

```

```

TYPE(I)=6RSTORED                                10700
S=S+1                                             10710
ITEMP1=CONVERT(S,2R L)                          10720
CALL OUTCDA2 (Z,ITEMP1,10H,                      10730
              //)
ITEMP=CODELOC(K)                                10740
K=K-1                                             10750
OUTCODE(ITEMP+2,3)=ITEMP1                       10760
ITEMP3=OPERAND(I)                              10770
146 CALL OUTCDA2 (Z,CONT(ITEMP3+2),10H.,        10780
               // )
ITEMP3=LINK(CONT(ITEMP3),1)                    10790
IF (ITEMP3.GT.0) GO TO 146                      10800
S=S+1                                             10810
ITEMP1=CONVERT(S,2R L)                          10820
CALL OUTCDA2 (Z,ITEMP1,10H,                      10830
              //)
OUTCODE(ITEMP,3)=ITEMP1                        10840
ITEMP=ITEMP+1                                   10850
OUTCODE(ITEMP,3)=CONVERT(1BYTE(CONT(OPERAND(I)),10),2R ) 10860
OUTCODE(ITEMP,5)=ITEMP1                        10870
P=P+1                                             10880
ITEMP=CONVERT(P,2R T)                          10890
CALL OUTCDA3 (Z,5H0 ,ITEMP,10H,                10900
              //)
S=S+1                                             10910
ITEMP1=CONVERT(S,2R L)                          10920
CALL OUTCDA2 (Z,ITEMP1,10H,                      10930
              //)
CALL LOOP82D (OPERAND(I),ITEMP1,OUTCODE)        10940
TYPE(I)=6RSTORED                                10950
OPERAND(I)=ITEMP                                10960
RETURN                                           10970
IF/PSI33A1/<SERIAL CL.>THEN/PSI82A2/<SER.CL.>ELSE/PSI82A3/<SER.CL.>FI 10980
IF/PSI33A1/<SERIAL CL.>THEN/PSI82A2/<SER.CL.>FI/PSI82B3/ 10990
                                                    11000
                                                    11010
                                                    11020
ENTRY PSI82A2
  IF(.NOT.SAMETYP(N(I),BOOL,REF))
1    CALL ERROR(110HIF STATEMENT WITH NON-BOOLEAN TEST.
2
  IF (.NOT.(TYPE(I).EQ.9RNONSTORED)) GO TO 147 11040
P=P+1                                             11050
ITEMP=CONVERT(P,2R T)                          11060
CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP,10H,        11070
              //)
147 CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z)        11080
I=I-1                                             11090
CALL OUTCDA4 (Z,CONVERT(P,2R T),3H=01,0,10H.,  11100
              //)
CODELOC(K)=Z                                    11110
;                                                 11120
; CODELOC CARRIES LOCATION TO INSERT FIXUP LABEL HERE. 11130
;                                                 11140
148 CALL BLOKEND(P,MODE)
  CALL BLOKBGN(P,MODE)
ITEMP=Z+3                                        11180
K=K+1                                            11190
CODELOC(K)=ITEMP                                11200
K=K+1                                            11210
CODELOC(K)=ITEMP                                11220
Z=ITEMP+1                                        11230

```

```

RETURN 11240
; 11250
ENTRY PSI82A3 11260
IF (OPFRAND(I).EQ.ITEMP) GO TO 149 11270
ITEMP=CONVERT(IBYTE(STORAGE(BLOKNUM),10)+1,2R T) 11280
CALL OUTCDA4 (2,OPERAND(I),1H,ITEMP,10H, 11290
TYPE(I)=6RSTORED 11300
OPERAND(I)=ITEMP 11310
149 CALL OUTCDA2 (2,0,10H., 11320
; 11330
; INSERT SKIP CLAUSE LABEL HERE. 11340
; 11350
K=K-1 11360
ITEMP1=CODELOC(K) 11370
CODELOC(K)=7 11380
S=S+1 11390
ITEMP=CONVERT(S,2R L) 11400
CALL OUTCDA2 (2,ITEMP,10H: 11410
OUTCODE(ITEMP1,3)=ITEMP 11420
GO TO 148 11430
; 11440
ENTRY PSI82A4 11450
CALL BLOKEND(P,MODE)
P=P+1
ITEMP=CONVERT(P,2R T) 11470
IF (OPERAND(I).EQ.ITEMP) GO TO 150 11480
CALL OUTCDA4 (2,OPERAND(I),1H,ITEMP,10H, 11490
; 11500
150 K=K-1 11510
; 11520
; BALANCING OF CONSEQUENCE AND ALTERNATIVE: 11530
; 11540
IF (.NOT.SAMETYP(N(I-1),N(I),REF)) GO TO 151 11540
IF (REF.EQ.0) GO TO 152 11550
GO TO 153 11560
151 IF (SAMETYP(N(I),N(I-1),REF)) GO TO 155 11570
CALL ERROR(110)THE CONSEQUENCE AND ALTERNATIVE OF AN IF STATEM
1ENT DO NOT HAVE THE SAME A POSTERIORI MODE.
; 11590
; NO BALANCING NEEDED. 11600
; 11610
152 S=S+1 11620
ITEMP1=CONVERT(S,2R L) 11630
CALL OUTCDA2 (2,ITEMP1,10H: 11640
ITEMP2=CODELOC(K) 11650
OUTCODE(ITEMP2,1)=ITEMP1 11660
K=K-1 11670
GO TO 154 11680
; 11690
; LOWER THE CONSEQUENCE TO THE LEVEL OF THE ALTERNATIVE,AND CAUSE THE A 11700
; TIVE TO SKIP OVER THE LOWERING CODE. 11710
; 11720
153 N(I-1)=N(I) 11730
CALL OUTCDA2 (2,0,10H., 11740
ITEMP=7 11750
S=S+1 11760

```

```

ITEMP1=CONVERT(S,2R L) 11770
ITEMP2=CODELOC(K) 11780
K=K-1 11790
CALL OUTCDA2 (Z,ITEMP1,10H: 11800
OUTCODE(ITEMP2,1)=ITEMP1 11810
CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z) 11820
S=S+1 11830
ITEMP1=CONVERT(S,2R L) 11840
CALL OUTCDA2 (Z,ITEMP1,10H: 11850
OUTCODE(ITEMP,1)=ITEMP1 11860
154 I=I-1
RETURN 11890
; 11900
; LOWER THE ALTERNATIVE, THEN INSERT JUMP CODE FOR THE CONSEQUENCE. 11910
; 11920
155 CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z) 11930
S=S+1 11940
ITEMP1=CONVERT(S,2R L) 11950
CALL OUTCDA (Z,ITEMP1,10H: 11960
ITEMP2=CODELOC(K) 11970
K=K-1 11980
OUTCODE(ITEMP2,1)=ITEMP1 11990
GO TO 154 12000
; 12010
; ENTRY PSI82B3 12020
; 12030
; INSERT ALTERNATIVE CODE AND PROVIDE A JUMP TO AVOID THIS CODE. 12040
; 12050
ITEMP=STORAGE(BLOKNUM) 12060
P=IRYTE(ITEMP,10)+1 12070
ITEMP1=CONVERT(P,2R T) 12080
CALL BLOKEND (0,0) 12090
CALL BLOK9GN (0,0) 12100
IF (OPERAND(I).EQ.ITEMP1) GO TO 156 12110
CALL OUTCDA4 (Z,OPERAND(I),1H:,ITEMP1,10H, 12120
TYPE(I)=6RSTORED 12130
OPERAND(I)=ITEMP1 12140
156 CALL OUTCDA2 (Z,0,10H:, 12150
S=S+1 12160
ITEMP2=CONVERT(S,2R L) 12170
K=K-1 12180
ITEMP3=CODELOC(K) 12190
K=K-1 12200
CALL OUTCDA2 (Z,ITEMP2,10H: 12210
OUTCODE(ITEMP3,3)=ITEMP2 12220
CALL OUTCDA3 (7,5H0 ,ITEMP1,10H, 12230
S=S+1 12240
ITEMP2=CONVERT(S,2R L) 12250
OUTCODE(7-2,1)=ITEMP2 12260
CALL OUTCDA2 (Z,ITEMP2,10H: 12270
RETURN 12280
ENTRY PSI84A 12290
C 12300
C TEST TO SEE WHETHER WE ARE COMPILING THE FIRST CLAUSE IN A ROW OF CL 12310
C THEN,IF TRUE,INSERT A JUMP AROUND THE ROW OF CLAUSE TO WHERE THE EVAL 12320

```

```

JUMP TABLE BEGINS.
P=IBYTE(STORAGE(BLOCKNUM),10)
ITEMP=CONVERT(1+P,2R T)
IF (OPERAND(I).EQ.ITEMP) GO TO 157
CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP,10H, //)
157 CALL OUTCDA3 (Z,10HRETURN RJ,8H(30*47),,10HRJ., //)
ITEMP1=IFETCH(4)
IF (CODELOC(K).EQ.-1) GO TO 158

THIS TEST DISTINGUISHES THE FIRST CLAUSE OF THE ROW.

K=K+1
CODELOC(K)=-1
S=S+1
ITEMP3=CONVERT(S,2R L)

THE DOUBLE USE OF <ROW OF CL.> WITHIN CASE STATEMENTS AND COLLATERAL
REQUIRES SOME BLANKING OF OUTCODE CARDS PUT IN BY PSI33A1 DURING PSIA

ITEMP2=CODELOC(K-2)
CALL OUTCDA2 (ITEMP2,ITEMP3,10H: //)
CODELOC(K-1)=ITEMP3
CODELOC(K-2)=ITEMP2-1
158 CALL STRIND (CODELOC(K-1),ITEMP1+1)
OPERAND(I)=ITEMP1

STORED ON OPERAND(I) IN A LIST CELL.

S=S+1
ITEMP3=CONVERT(S,2R L)
CODELOC(K-1)=ITEMP3
CALL OUTCDA2 (Z,ITEMP3,10H: //)
RETURN

THE LABEL AT WHICH CODE BEGINS FOR EACH ROW IN A <ROW OF CLAUSE> IS

ENTRY PSIA3A

DO PSIA4A WORK FIRST, THEN PIECE TOGETHER THE COMPILE-TIME REPRESENTAT
OF <ROW OF CLAUSE>.

P=IBYTE(STORAGE(BLOCKNUM),10)
ITEMP=CONVERT(P+1,2R T)
IF (OPERAND(I).EQ.ITEMP) GO TO 159
CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP,10H, //)
159 CALL OUTCDA3 (Z,10HRETURN RJ,8H(30*47),,10HRJ., //)
ITEMP=IFETCH(4)
CALL STRIND (CODELOC(K-1),ITEMP+1)

C CONSTRUCT LIST OF <ROW OF CL.> PROCEDURES.
C
CALL STRIND (ITEMP,OPERAND(I-1))
C
C CONSTRUCT A COMPILE TIME STRUCTURE SKELETON REPRESENTATION OF THE DATA

```

```

ITEMP=IFETCH(1)
CALL SETIND (N(I),ITEMP,2)
I=I-1
ITEMP1=IFETCH(1)
CALL SETIND (N(I),ITEMP1,2)
CALL SETIND (ITEMP,ITEMP1,1)
CALL PUTIND (2,ITEMP1,10)
N(I)=ITEMP1
RETURN

ENTRY PSI83B

ITEMP=IFETCH(1)
CALL SETIND (N(I),ITEMP,1)
I=I-1
CALL SETIND (I(I),ITEMP,2)
N(I)=ITEMP
CALL PUTIND (1+IBYTE(CONT(N(I+1))),10),ITEMP,10)
CALL SETIND (OPERAND(I+1),OPERAND(I))
RETURN
ENTRY PSI85A2

INSERT JUMP AROUND <POW OF CLAUSE> CODE, THEN DO BLOCKEND ROUTINE.

K=K-2
ITEMP=CODELOC(K)-2
ITEMP1=ITEMP+1
DO 160 I1=ITEMP,ITEMP1
160 OUTCODE(I1,1)=2H//
CALL OUTCODE2 (CODELOC(K),CODELOC(K+1),10H., //)
GO TO 184

ENTRY PSI48B2

AT THIS POINT, THE MODE OF THE <VIRT. MODE DECLARER> MUST BE EITHER S
OR ARRAY, WITH NO PROC OR REF PREFIXES. IN ADDITION, THE <COLLATERAL
AND ITS A POSTERIORI MODE MUST BE THE SAME LENGTH.

ITEMP1=CONT(N(I))
ITEMP2=CONT(N(I-1))
P=P+1
ITEMP3=IBYTE(ITEMP1,10)
IF (ITEMP3.EQ.IBYTE(ITEMP2,10)) GO TO 161
CALL ERROR (110H THE <COLLATERAL CLAUSE> AND ITS A POSTERIORI MODE AR
1E NOT THE SAME LENGTH.
RETURN
161 ITEM4=LINK(ITEMP1,3)
IF (ITEM4.EQ.3RARY) GO TO 162
IF (ITEM4.EQ.3RSTR) GO TO 164
CALL ERROR (110H THE A POSTERIORI MODE OF A <MODE CAST> WITH A <CO
1LLATEPAL CLAUSE> IS NOT ARRAY OR STRUCTURE, AS IT MUST BE. )
RETURN

AT THIS POINT, WE FIND THE LOWEST MODE IN THE <ROW OF CLAUSE>, CONVERT

```



```

CIRCULAR LIST RESEMBLING A STRUCTURE, AND CALL THE LOWER SUBROUTINE. 13430
152 CALL OUTCDA4 (Z,7HSAVEU, ,LHCURR,10H+1U,SAVEU,10H, 13440
13450
13460
CREATE A RUN-TIME ARRAY:STEP 1: GENERATE THE DESCRIPTOR. 13470
13480
ITEMP4=CONVERT(ITEMP3+2,2R ) 13490
CALL OUTCDA6 (Z,LHCURR,1H+,ITEMP4,1H,LHCURR,10H, 13500
CALL OUTCDA3 (Z,10H0100000000,10H0000 U,10H, 13510
CALL OUTCDA3 (Z,10HU+2 U,10H(0+17),U+1,10H+U, 13520
CALL OUTCDA3 (Z,10H0300000100,10H0000000001,10H+U, 13530
CALL OUTCDA (Z,CONVERT(ITEMP3,2R ),10H+U(18+35,8H), 13540
13550
STEP 2: CALL LOWER, AFTER GENERATING CIRCULAR LIST FOR LOWEST MODE. 13560
13570
ITEMP4=MINIMUM(N(I)) 13580
ITEMP5=IFETCH(1) 13590
CALL SETIND (ITEMP5,ITEMP5,1) 13600
CALL SETIND (ITEMP4,ITEMP5,2) 13610
CALL LOWER (ITEMP3,ITEMP5,N(I),OPERAND(I),Z,P) 13620
13630
STEP 3: COMPLETE WORK ON THE RUN-TIME DATA ELEMENT. 13640
13650
153 ITEMP4=CONVERT(P,2R T) 13660
CALL OUTCDA3 (Z,10HPOPU, U ,ITEMP4,10H,POPU, 13670
I=I-1 13680
OPERAND(I)=ITEMP3 13690
RETURN 13700
13710
AT THIS POINT, A STRUCTURE IS GENERATED AT RUN TIME FROM THE ELEMENTS 13720
<COLLATERAL CLAUSE>. 13730
13740
164 CALL OUTCDA4 (Z,6HSAVEU, ,LHCURR,10H+1 U,SAVE,10HU, 13750
CALL OUTCDA6 (Z,LHCURR,1H+,CONVERT(ITEMP3,2R ),1H,LHCURR,10H, 13760
1 13770
CALL LOWER (ITEMP3,LINK(ITEMP2,2),OPERAND(I),Z,P) 13780
GO TO 163 13790
13800
13810
ENTRY PSI49A1 13820
13830
THE MODE ASSOCIATED ON THE N-STACK WITH A <VOID CAST> IS 0. NOTE THA 13840
THE VALUE OF P DOES NOT CHANGE HERE BECAUSE OF THE ALGOL 68 CONVENTI 13850
OF LEAVING THE DESTINATION OF AN ASSIGNATION AS ITS VALUE AND BECAUSE 13860
OUR PSI33A2 RESETS P. 13870
13880
N(I)=0 13890
OPERAND(I)=0 13900
RETURN 13910
13920
ENTRY PSI48A2 13930
13940
P = P + 1 13950
13960
ITEMP=CONVERT(P,2R T) 13970

```


	RETURN	14910
C		14920
	ENTRY PSI15B1	14930
C		14940
	K=K+1	14950
	CODELOC(K)=Z	14960
	RETURN	14970
C		14980
	ENTRY PSI16C	14990
C		15000
	ITEMP=IFETCH(1)	15010
	CALL SETIND (N(I),ITEMP,2)	15020
	CALL SETIND (N(I-1),ITEMP,1)	15030
	CALL PUTIND (IBYTE(CONT(N(I-1)),10),N(I),10)	15040
	I=I-1	15050
	N(I)=ITEMP	15060
	RETURN	15070
C		15080
	ENTRY PSI77A1	15090
C		15100
C	IN THE PROCEDURE CALL, THE <PRIMARY> MUST BE LOWERED TO ITS FIRST MOD	15110
C	A PROCEDURE WITH PARAMETERS.	15120
C		15130
	NI=N(I)	15140
	ITEMP3=LOGF(ITEMP4)	15150
	IF (NI.GT.0) GO TO 173	15160
	172 CALL ERROR (110HPROCEDURE CALL ATTEMPTED ON A NON-PROCEDURE <PRIMA	15170
	RY>.	15180
	RETURN	15190
	173 ITEMP=CONT(NI)	15200
	ITEMP1=LINK(ITEMP,3)	15210
C		15220
C	THE MODE PREFIXES MUST BE SOME SEQUENCE OF REF, STOPPING AT THE FIRST	15230
C	WITH PARAMETERS.	15240
C		15250
	IF (ITEMP1.NE.3RREF) GO TO 175	15260
	ITEMP2=3RREF	15270
	174 ITEMP5=IFETCH(1)	15280
	CALL SETIND (ITEMP2,ITEMP5,3)	15290
	CALL SETIND (ITEMP5,ITEMP3,2)	15300
	ITEMP3=ITEMP5	15310
	NI=LINK(ITEMP,2)	15320
	GO TO 173	15330
	175 IF (ITEMP1.NE.3RPROC) GO TO 172	15340
	IF (LINK(ITEMP,2).GT.0) GO TO 176	15350
	ITEMP2=3RPROC	15360
	GO TO 174	15370
	176 REF=ITEMP4	15380
	P=P+1	15390
	ITEMP4=CONVERT(P,2R T)	15400
	IF (OPERAND(I).EQ.ITEMP4) GO TO 177	15410
	CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP4,10H, //	15420
	OPERAND(I)=ITEMP4	15430
	TYPE(I)=6RSTORED	15440
	177 N(I)=NI	15450

```

      CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z) 15460
3 . 15470
3 NEXT, STORE A LABEL ON N(I) FOR USE IN SETTING UP THE ACTUAL PARAMET 15480
3 LIST, AND PROVIDE A JUMP AROUND THE ACTUAL PARAMS. LIST FOR LATER CONF 15490
3 15500
      I=I+1 15510
      S=S+1 15520
      ITEMP4=CONVERT(S,2R L) 15530
      N(I)=ITEMP4 15540
      CALL OUTCDA4 (Z,0,0,ITEMP4,10H; 15550
      TYPE(I)=Z 15560
      CALL SETDIR (P+1,TYPE(I),2) 15570
      RETURN 15580
C 15590
      ENTRY PSI78B1 15600
C 15610
C THIS RULE CREATES A SEQUENCE OF LIST CELLS CONTAINING MODE INFORMATIO 15620
C RUN-TIME LABELS LOCATING WHERE EACH <UNITARY CLAUSE> BEGINS. 15630
C 15640
178 ITEMP=LINK(TYPE(I-1),2) 15650
      ITEMP3=50 15660
      ITEMP1=CONVERT(ITEMP,2R T) 15670
      IF (ITEMP1.EQ.OPERAND(I)) GO TO 179 15680
      CALL OUTCDA4 (Z,OPERAND(I),1H,ITEMP1,10H, 15690
179 P=ITEMP-1 15700
      CALL OUTCDA3 (Z,10HRETURN, RJ,10H(30,47), ,10HRJ., 15710
      ITEMP2=IFETCH(2) 15720
      CALL STRIND (N(I-1),ITEMP2+1) 15730
C 15740
C STORES LABEL IN SECOND WORD. 15750
C 15760
      CALL SETINC (N(I),ITEMP2,2) 15770
C 15780
C STORES MODE LINK IN FIRST WORD. 15790
C 15800
      OPERAND(I)=ITEMP2 15810
      IF (ITEMP3.EQ.0) RETURN 15820
      S=S+1 15830
      ITEMP2=CONVERT(S,2R L) 15840
      CALL OUTCDA2 (Z,ITEMP2,10H; 15850
      N(I)=ITEMP2 15860
      CALL SETDIR (ITEMP,TYPE(I),2) 15870
      RETURN 15880
C 15890
      ENTRY PSI78A1 15900
C 15910
C ABBREVIATED VERSION OF PSI78B1 WITHOUT INSERTION OF THE NEXT LABEL IN 15920
C OUTPUT CODE. 15930
C 15940
      ITEMP3=0 15950
      GO TO 178 15960
C 15970
      ENTRY PSI78B2 15980
C 15990
C LINKS TOGETHER THE LIST OF LABELS AND MODE LINKS. 16000

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> CALL SETIND (OPERAND(I), OPERAND(I-1), 1) 16010
> I=I-1 16020
> RETURN 16030
> 16040
> ENTRY PSI77A2 16050
> 16060
> 16070
> OPERAND(J) CONTAINS THE LIST OF ENTRY LABELS AND DATA TYPES FOR EACH 16080
> ACTUAL PARAMETERS. P WAS RESET TO ITS RETURN LEVEL - 1 BY PSI78A1. LH 16090
> SAVED BEFORE PARAMETER EVALUATION, SO THAT ITS INITIAL VALUE BEFORE ST 16100
> PARAMETERS ON IT CAN BE RETRIEVED TO PASS TO THE PROCDEF. AT END OF 16110
> EVALUATION, CHECK FOR SAME LENGTH OF P-LISTS. TYPE(I-1) CONTAINS TH 16120
> TO INSERT A JUMP OVER THE PARAMETERS, AND N(I-2) CONTAINS THE A POSTE 16130
> PROCEDURE MODE. 16140
> 16150
> ITEMP=TYPE(I-1).AND.777777B 16160
> S=S+1 16170
> ITEMP1=CONVERT(S, 2R L) 16180
> OUTCODE(ITEMP, 1)=ITEMP1 16190
> OUTCODE(ITEMP, 2)=10H., 16200
> CALL OUTCOA2 (Z, ITEMP1, 10H., //) 16210
> P=P+1 16220
> NI=CONVERT(P, 2R T) 16230
> CALL OUTCOA4 (Z, 7HSAVE U, LHCURR, 5H+1U, 10HSAVE U, //) 16240
> 16250
C LOWER<ACTUAL PARAMETERS> TO A POSTERIORI LEVEL, THEN STORE IN TURN IN 16260
C 16270
C ITEMP5=0 16280
C 16290
C COUNT OF PARAMETERS 16300
C 16310
C ITEMP1=OPERAND(I) 16320
C ITEMP3=LINK(CONT(N(I-2)), 1) 16330
C 16340
C ITEMP1 AND ITEMP3 ARE, RESPECTIVELY, LINKS TO THE A PRIORI AND A POSTE 16350
C PARAMETER LISTS. 16360
C 16370
180 ITEMP2=CONT(ITEMP1) 16380
ITEMP4=CONT(ITEMP3) 16390
ITEMP5=1+ITEMP5 16400
IF (SAMETYP(LINK(ITEMP2, 2), LINK(ITEMP4, 2), REF)) GO TO 181 16410
CALL ERROR(110H THE PARAMETER IN THE JTH POSITION OF THE PROCED
180 CALL DOES NOT MATCH ITS A POSTERIORI MODE. J IS... )
CALL ERROR1(CONVERT(ITEMP5, 2R ))
GO TO 180
181 CALL OUTCOA6 (Z, 10HSAVE RETUR, 10HN, J+2*RETU, 3HRN, CONT(ITEMP1+1), 1
10H., POP RETU, 10HRN, //)
CALL OUTEVAL (P, REF, OUTCODE, Z, 1000-Z)
CALL OUTCOA4 (Z, LHCURR, 3H+1U, LHCURR, 10H, //)
CALL OUTCOA4 (Z, NI, 1H., LHCURR, 10H, //)
ITEMP1=LINK(ITEMP2, 1)
ITEMP3=LINK(ITEMP4, 1)
IF ((ITEMP1.GT.0).AND.(ITEMP3.GT.0)) GO TO 180
IF (ITEMP1.EQ.ITEMP3) GO TO 182
CALL ERROR (110H ILLEGAL PROCEDURE CALL--THE NUMBER OF ACTUAL AND F

```



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ENTRY PSI73A2                                     17130
C                                                     17140
ITEMP=IFETCH(1)                                    17150
CALL SETIND (3RPRC,ITEMP,3)                        17160
CALL SETIND (N(I),ITEMP,2)                        17170
CALL OUTCDA4 (Z,OPERAND(I),10H * VALUE, ,10H[RETURN]., ,10H 17180
1//)                                                17190
S=S+1                                              17210
ITEMP1=CONVERT(S,2R L)                             17220
I=I-1                                              17230
ITEMP2=N(I)                                         17240
OUTCODE(ITEMP2,1)=ITEMP1                          17250
CALL OUTCDA2 (Z,ITEMP1,10H)                        17260
N(I)=ITEMP                                         17270
CALL PRCEAD(P,0)
RETURN                                             17300
C                                                     17310
C *****                                           17320
C                                                     17330
C                                                     17390
C                                                     17400
END
SUBROUTINE OUTEVAL (P,REF,OUTCODE,Z,MAX)           10
INTEGER P,REF,OUTCODE,Z,CONT,CONVERT              20
DIMENSION OUTCODE(MAX,6)                          30
ITEMP=CONVERT(P,2R T)                             40
101 ITEM3=LINK(CONT(REF),2)                        50
C                                                     60
C REF POINTS TO THE LIST CREATED BY SAMETYP CONTAINING INFORMATION 70
C ABOUT THE CORRECT SEQUENCE FOR DEPROCEDURING AND DEREFERENCING 80
C THE EXPRESSION CURRENTLY BEING EVALUATED.        90
C                                                     100
IF (ITEM3.EQ.0) RETURN                             110
IF (ITEM3.EQ.3RREF) GO TO 102                      120
C                                                     130
C THE REMAINING ALTERNATIVE, WHICH MUST BE TRUE, IS THAT ITEM SIGNALS A 140
C PROCEDURE CALL. THE NUMBER OF PARAMETERS IS FOUND IN LINK 3 OF THE 150
C WORD POINTED TO BY REF.                          160
C                                                     170
ITEM4=CONVERT(LINK(CONT(REF),3),2R )              180
CALL OUTCDA4 (Z,ITEM4,10H*NUMBER OF,10HPARAMATERS,8H, //) 190
P=P-LINK(CONT(REF),3)                             200
ITEMP=CONVERT(P,2R T)                             210
CALL OUTCDA4 (Z,10HSAVE RETUR,3HN, ,10HITEMP*RJ( ,10H30*47), //) 220
C                                                     230
C CONTROL IS PASSED TO THE PROCEDURE9              240
C                                                     250
CALL OUTCDA2 (Z,10HJ+2*RETURN,10H, RJ., //)        260
CALL OUTCDA5 (Z,10HPOPRETURN,,10HVALUE RETU,3HRN*,ITEMP,10H, 270
1 //)                                               280
CALL OUTCDA2 (Z,10HPOP VALUE ,10HRETURN, //)        290
REF=LINK(CONT(REF),1)                             300
GO TO 101                                           310
C                                                     320
C THE NUMBER OF REFERENCE LEVELS IS FOUND IN LINK 3 OF THE WORD POINTED 330
C TO BY REF.                                       340

```

```

C
102 ITEM4=LINK(CONT(REF),3)
GO TO 104
103 ITEM4=ITEM4-1
104 CALL OUTCDA5 (Z,1H[,ITEMP,2H],ITEMP,10H,      //)
IF (ITEM4.GT.0) GO TO 103
REF=LINK(CONT(REF),1)
GO TO 101
C
END
INTEGER FUNCTIONMINIMUM(NI)
LOGICAL SAMETYP
INTEGER REF,CONT
ITEMP1=CONT(NI)
LINK1=LINK(ITEMP1,2)
ITEMP=LINK(ITEMP1,1)
101 ITEM1=CONT(ITEMP)
LINK2=LINK(ITEMP1,2)
IF (SAMETYP(LINK1,LINK2,REF)) GO TO 102
C
C AT 5, LINK2 IS THE LOWER DATA TYPE.
C
IF (SAMETYP(LINK2,LINK1,REF)) GO TO 103
CALL ERROR (82)
RETURN
102 LINK1=LINK2
103 ITEM=LINK(ITEMP1,1)
IF (ITEM.GT.0) GO TO 101
MINIMUM=LINK1
RETURN
C
END
SUBROUTINE BALANCE (LINK1,ITEMP1,ITEMP2,S,P0,Z,OUTCODE)
DIMENSION OUTCODE(5000,6)
LOGICAL SAMETYP
INTEGER S,P,Z,OUTCODE,CONT,P0
ITEMP3=LINK(CONT(ITEMP2),2)
101 S=S+1
ITEMP=CONVERT(S,2R L)
P=P0+1
CALL STRIND (ITEMP,ITEMP1+2)
CALL OUTCDA2 (Z,ITEMP,10H,      //)
CALL OUTCDA3 (Z,10HSAVE RETUR,10HN,J+2+RETU,10RN,      // )
CALL OUTCDA2 (Z,CONT(ITEMP1+1),10H.,      //)
CALL OUTEVAL (P,REF,OUTCODE,Z,1000-Z)
CALL STRIND (Z+1,ITEMP1+3)
C
C HERE,THE ESCAPE ADDRESS IS PUT INTO THE THIRD LIST ELT. DATUM.
C SEE PS18203 FOR DETAILS. NOTE THAT THE ROWS OF CLAUSE USE TICODELOC(K
C AS THE VALUE RETURN ABOVE.
C
CALL OUTCDA2 (Z,0,10H.,      //)
ITEMP1=CONT(ITEMP1)
ITEMP2=LINK(CONT(ITEMP2),1)
IF (ITEMP1.GT.0) GO TO 101

```


1, STORAGE(50), PSTORE(50)	40
1, INTEGER HASH1, HASH2, STORAGE, BLOKNUM, PRLEVEL, PRCOUNT, BLCOUNT, PSTORE	50
1, CONVERT	60
LOGICAL SAPETYP	70
;	80
ENTRY INITIAL	90
;	100
;	110
;	120
PRLEVEL=0	130
BLOKNUM=0	140
BLCOUNT=0	150
PRCOUNT=0	160
MBCOUNT=0	
MPCOUNT=0	
L=IFETCH1(0)	170
;	180
;	190
INITIALIZES ALL COMPILE TIME LISTS	200
;	210
RETURN	220
;	230
ENTRY FINAL	240
;	250
;	260
CALLLED WHEN COMPILER REQUESTS TALLY OF NAMES IN NAMES AND PROCS TABLE	270
;	280
NAME=MBCOUNT	290
;	300
MAXIMUM NAMES IN STATIC ENVIRONMENT	
;	320
MODE=MPCOUNT	330
;	340
MAXIMUM NO. OF NAMES IN PROCEDURE ENVIRONMENT.	350
;	360
RETURN	370
;	380
ENTRY BLOKEGN	390
;	400
BLOKNUM=BLOKNUM+1	
STORAGE(BLOKNUM)=BLCOUNT	
CALL SETDIR(PRCOUNT, STORAGE(BLOKNUM), 2)	
CALL SETDIR(NAME, STORAGE(BLOKNUM), 3)	
101 L=IFETCH2(0)	410
;	420
;	430
COMPILE-TIME LISTS BOOKKEEPING	440
;	450
RETURN	460
;	470
ENTRY BLOKEND	480
;	490
L=BLCOUNT-STORAGE(BLOKNUM)-1	500
IF (L.LT.0) GO TO 103	510
DO 102 I=0,L	520
J=NAMES1(BLCOUNT-I,1)	530
102 HASH1(J,1)=0	
103 BLCOUNT=STORAGE(BLOKNUM).AND.77777B	

```

        PRCOUNT=LINK(STORAGE(BLOKNUM),2)
        NAME=LINK(STORAGE(BLOKNUM),3)
        BLOKNUM=BLOKNUM-1
104 L=IFETCH3(0)
CALLS THE LIST GARBAGE COLLECTER.
RETURN
ENTRY PROCBGN
PRLEVEL=PRLEVEL+1
PSTORE(PLEVEL)=PCOUNT
CALL SETDIR(NAME,PSTORE(PLEVEL),2)
GO TO 101
ENTRY PROCEND
DELETES COMPILE TIME NAMES IN CURRENT PROCEDURE BLOCK, BUT LEAVES THE
TIME LOCATIONS INTACT.
L=PCOUNT-1-LINK(PSTORE(PLEVEL),1)
NAME=LINK(PSTORE(PLEVEL),2)
IF (L.LT.0) GO TO 106
DO 105 I=0,L
J=NAMES2(PCOUNT-I,1)
105 HASH2(J,1)=0
106 PRLEVEL=PRLEVEL-1
GO TO 104
ENTRY IHASHA
LOOK UP NAME AND MODE,OR LABEL AND LIST OF LOCATIONS TO BE FILLED,AND
RUN-TIME ALIAS PRECEDED BY ~.
LASTUSE=0
I=NAME
IF (PRLEVEL.EQ.0) GO TO 111
DO 110 L=0,196,1
I=MOD(I+47*L,197)+1
IA=HASH2(I,1)
IF (IA.NE.0) GO TO 108
IF (LASTUSE.GT.0) GO TO 107
GO TO 111
107 IHASHA=CONVERT(LASTUSE,2R~P)
MODE=LOCN(NAMES2(LASTUSE,2))
RETURN
108 IF (IA.NE.NAME) GO TO 109
LASTUSE=HASH2(I,2)
109 CONTINUE
110 CONTINUE
IF(LASTUSE.GT.0)GO TO 107
C NAME IS NOT WITHIN PROCEDURE ENVIRONMENT.
111 DO 115 L=0,196,1

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```

        I=MOD(I+47*L,197)+1
        IA=HASH1(I,1)
        IF (IA.NE.0) GO TO 113
        IF (LASTUSE.GT.0) GO TO 112
        CALL ERROR(110)THE TAG BELOW IS USED IMPROPERLY9 IT HAS NOT
        BEEN DECLARED IN THIS OR ANY GLOBAL BLOCKS9
        CALL ERROR1(NAME)
112     IHASHA=CONVERT(LASTUSE,2R~V)
        MODE=LOCF(NAMES1(LASTUSE,2))
        RETURN
113     IF (IA.NE.NAME) GO TO 114
        LASTUSE=HASH1(I,2)
114     CONTINUE
115     CONTINUE
        IF (LASTUSE.GT.0) GO TO 112
;
        ENTRY IHASHB
;
; ENTER NAME AND MODE ON NAMES TABLE, AND RETURN RUN-TIME ALIAS PREFIXED
;
        I=NAME
        IF (PRLEVEL.EQ.0) GO TO 118
        DO 116 L=0,196,1
            I=MOD(I+47*L,197)+1
            IA=HASH2(I,1)
            IF (IA.EQ.0) GO TO 117
;
; ENTER THE NAME.
;
116     CONTINUE
        CALL ERROR (110)HCOMPIL-TIME NAME TABLE OVERFLOW DURING A PROCEOUR
        IE DECLARATION.
        RETURN
117     PRCOUNT=PRCOUNT+1
        IF (PRCOUNT.GT.MPCOUNT) MPCOUNT=PRCOUNT
        HASH2(I,1)=NAME
        HASH2(I,2)=PRCOUNT
        NAMES2(PRCOUNT,1)=I
        NAMES2(PRCOUNT,2)=MODE
        MODE=LOCF(NAMES(PRCOUNT,2))
C
C FOR USE BY LABELS.
C
        IHASHB=CONVERT(PRCOUNT,2R~P)
        RETURN
118     DO 119 L=0,196,1
            I=MOD(I+47*L,197)+1
            IA=HASH1(I,1)
            IF (IA.EQ.0) GO TO 120
119     CONTINUE
        CALL ERROR (110)HCOMPIL-TIME NAME TABLE OVERFLOW OUTSIDE OF A PRO
        1CEDURE DECLARATION.
        RETURN
120     BLCOUNT=BLCOUNT+1
        IF (BLCOUNT.GT.MBCOUNT) MBCOUNT=BLCOUNT

```

```

HASH1(I,1)=NAME                                1570
HASH1(I,2)=BLCOUNT                            1580
NAMES1(BLCOUNT,1)=I                            1590
NAMES1(BLCOUNT,2)=MODE                         1600
MODE=LOCN(NAMES1(BLCOUNT,2))                  1610
THASHP=CONVERT(BLCOUNT,2R)V                   1620
RETURN                                          1630
                                           1640
ENTRY IHASHC                                   1650
                                           1660
THIS FUNCTION LOOKS UP OPERATORS IN THE NAME TABLE. WE ASSUME THAT AN
OPERATOR NOT PRESENT ON THE TABLE IS EITHER INVALID OR WIRED IN TO THE
RUN-TIME PILOT SYSTEM.                       1670
                                           1680
                                           1690
                                           1700
LASTUSE=0                                       1710
I=NAME                                          1720
IF (PPLEVEL,0,0) GO TO 124                    1730
DO 123 L=0,196,1                              1740
    I=MOD(I+47*L,197)+1                       1750
    IA=HASH2(I,1)                             1760
    IF (IA.NE.0) GO TO 121                     1770
    IF (LASTUSE.GT.0) GO TO 107                1780
    GO TO 124                                  1790
121    IR=HASH2(I,2)                           1800
                                           1810
; COMPARE OPERAND MODES.                      1820
;                                           1830
    IB=LINK(CONT(NAMES2(IB,2)),1)              1840
    IF ((IA.NE.NAME).AND.(.NOT.SAMETYP(MODE,IB,REF))) GO TO 122 1850
    LASTUSE=HASH2(I,2)                        1860
122    CONTINUE                               1870
123    CONTINUE                               1880
    IF (LASTUSE.GT.0) GO TO 107                1890
124 DO 126 L=0,196,1                          1900
    I=MOD(I+47*L,197)+1                       1910
    IA=HASH1(I,1)                             1920
    IF (IA.NE.0) GO TO 125                     1930
    IF (LASTUSE.GT.0) GO TO 112                1940
    IHASHC=0                                  1950
    RETURN                                    1960
125    IR=HASH1(I,2)                           1970
    IB=LINK(CONT(NAMES1(IB,2)),1)              1980
    IF ((IA.NE.NAME).AND.(.NOT.SAMETYP(MODE,IB,REF))) GO TO 126 1990
    LASTUSE=HASH1(I,2)                        2000
126    CONTINUE                               2010
C                                           2020
    IF (LASTUSE.GT.0) GO TO 112
    RETURN
END

```